



Bidirectional Search-Based Synthesis

Woosuk Lee

CSE9116 SPRING 2024

Hanyang University

Top-Down vs. Bottom-Up

- **Top-down search**

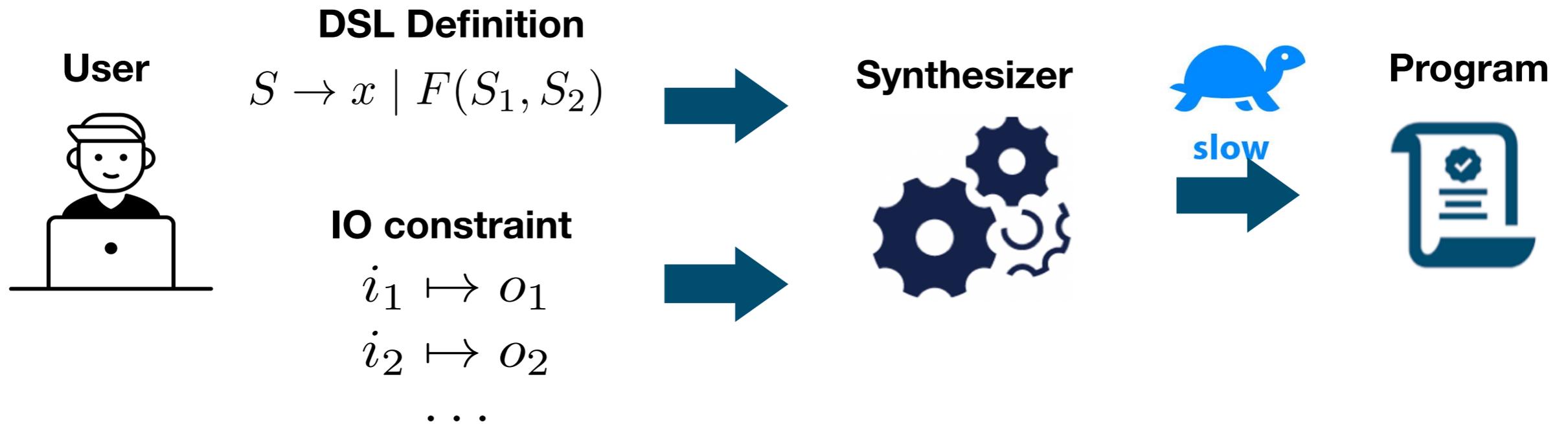
- Candidates are whole but might be incomplete (w/ holes); cannot always run on inputs but always relate to outputs
- Optimization: *top-down propagation*
 - Efficient but only applicable for specific kinds of languages

- **Bottom-up search**

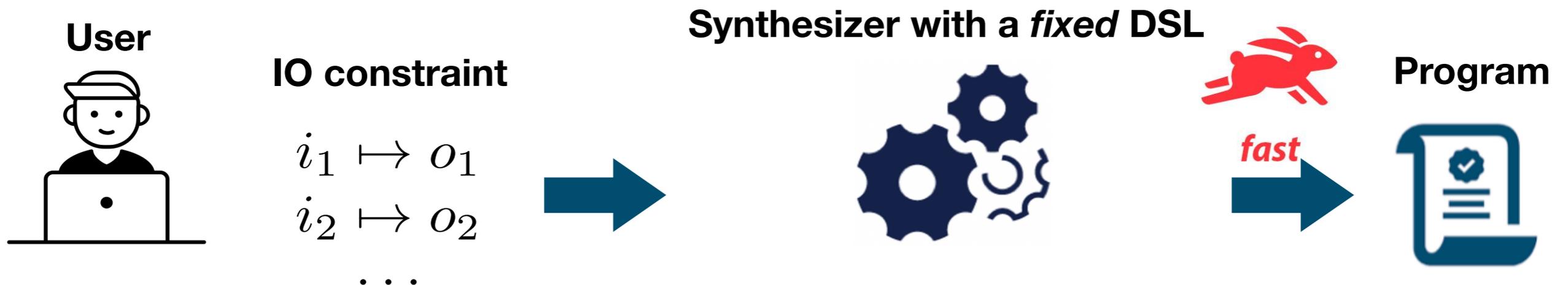
- Candidates are complete programs but might not be whole; can always run on inputs but cannot always relate to outputs
- Optimization: *observational equivalence reduction*
 - Generally applicable but inefficient

Dichotomy : “General-purpose” vs. “Domain-specific”

General-purpose synthesizer

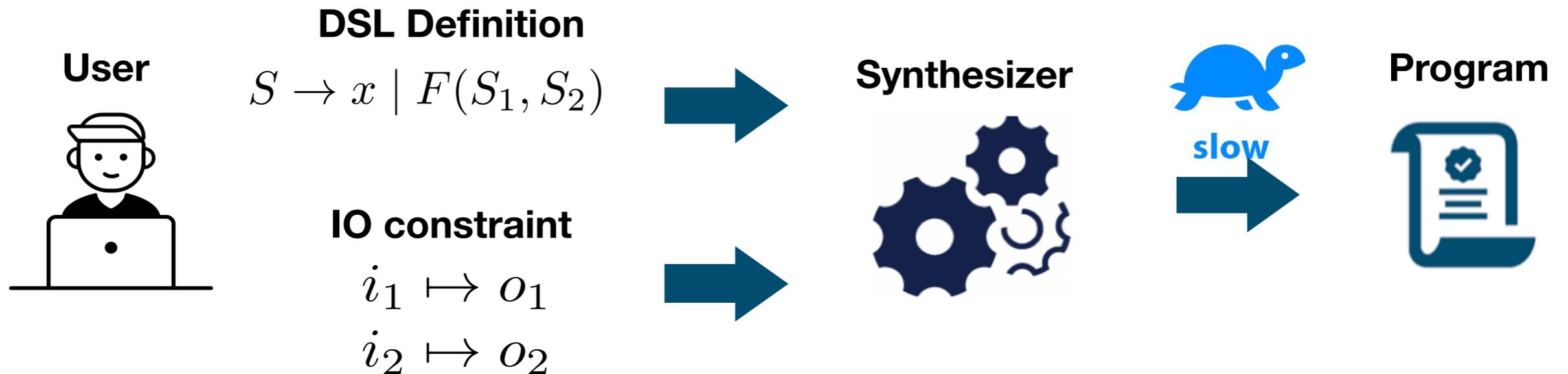


Domain-specific synthesizer



Dichotomy : “General-purpose” vs. “Domain-specific”

General-purpose synthesizer



- Encouraged by a standard formulation: Syntax-guided Synthesis (SyGuS)
- General search strategies (e.g., **Bottom-up enumeration**)

+ Broad application domains

– unscalable

Dichotomy : “General-purpose” vs. “Domain-specific”

- Domain-specific search strategies (e.g., **Top-down propagation**)

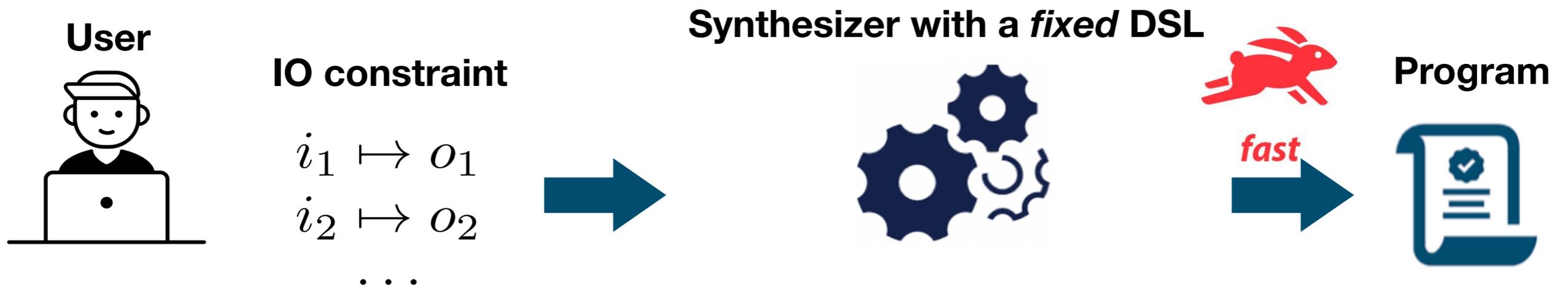
- Successful industrialization (e.g.,   )

+ very efficient

– only applicable for specific applications

...

Domain-specific synthesizer

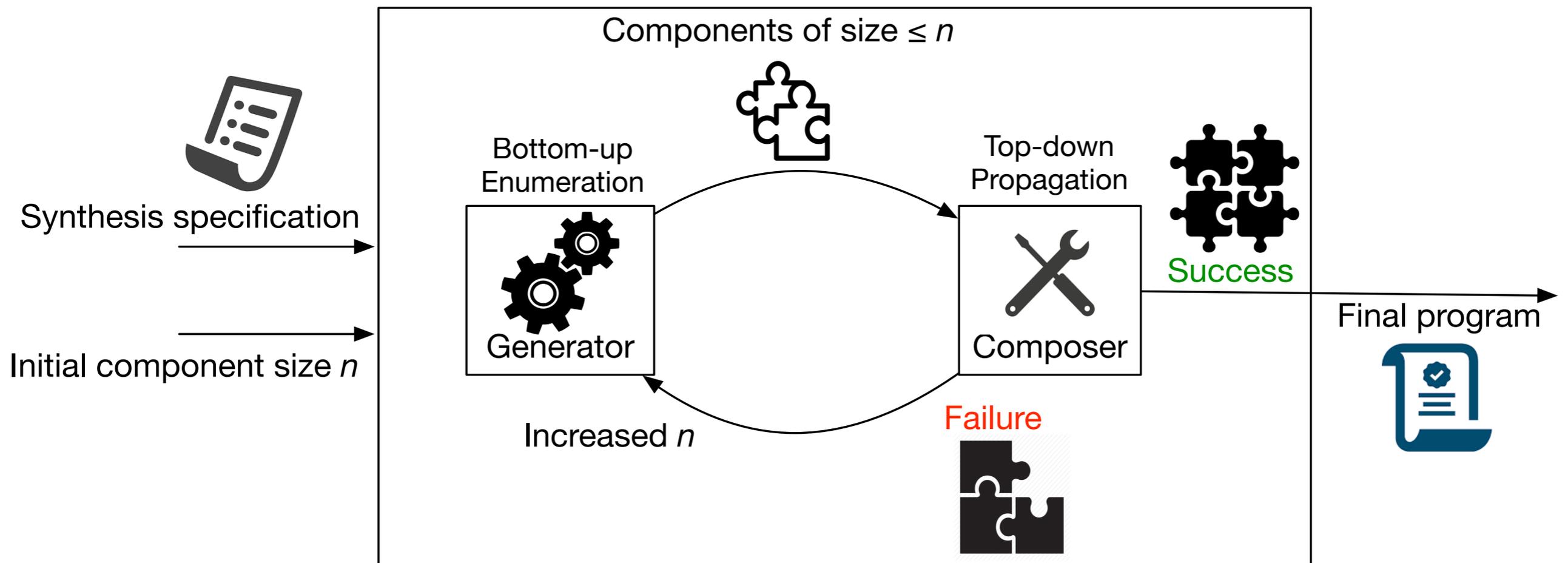


A Bidirectional Search Strategy

- Woosuk Lee, Combining the Top-down Propagation and Bottom-up Enumeration for Inductive Program Synthesis, POPL 2021
- Synergistically combine top-down and bottom-up search
- Generally applicable for a wide range of SyGuS instances
- <https://github.com/wslee/duet>

Overall Algorithm

Domain-**U**naware inductive synthesis via **E**numeration and **T**op-down propagation



Example

- Goal: function f converting a given phone number x into another format

-  Specification

Syntactic: $S \rightarrow x \mid \text{"+"} \mid \text{"_"} \mid \text{"-"} \mid \text{"."} \mid \text{ConCat}(S, S)$
| $\text{SubStr}(S, I, I)$
| $\text{IntToStr}(I)$
| $\text{Replace}(S, S, S)$

String concatenation

$I \rightarrow 1 \mid \dots \mid 9 \mid I - I \mid \text{Length}(S)$

Semantic:

$f(\text{"_ + 1_2.3"}) = \text{"1_2-3."} \wedge f(\text{"_ + 4_56.78"}) = \text{"4_56-78."}$

Example

- Goal: function f converting a given phone number x into another format

-  S

$\text{SubStr}(str, start_pos, len)$
e.g., $\text{SubStr}("abc", 0, 2) = "ab"$

Syntax: $S \rightarrow S \mid "-" \mid "." \mid \text{Concat}(S, S)$

| $\text{SubStr}(S, I, I)$

| $\text{IntToStr}(I)$

| $\text{Replace}(S, S, S)$

$I \rightarrow 1 \mid \dots \mid 9 \mid I - I \mid \text{Length}(S)$

Semantic:

$f("_{ } + 1_{ }2.3") = "1_{ }2-3." \wedge f("_{ } + 4_{ }56.78") = "4_{ }56-78."$

Example

- Goal: function f converting a given phone number x into another format

-  Specification

Syntactic: $S \rightarrow \text{Replace}(str, match, replacement) \quad S, S)$
|
| e.g., $\text{Replace}("aba", "a", "b") = "bba"$
|
| $\text{Replace}(S, S, S)$

$I \rightarrow 1 \mid \dots \mid 9 \mid I - I \mid \text{Length}(S)$

Semantic:

$f("_{_} + 1_{_}2.3") = "1_{_}2-3." \wedge f("_{_} + 4_{_}56.78") = "4_{_}56-78."$

Example

-  Solution: Size : 12 AST nodes

Replace(SubStr(Concat(x , ".") , 2, Length(x) - 1), ".", "-")

outputs: $\langle \text{"_+1_2.3."}, \text{"_+4_56.78."} \rangle$

outputs: $\langle \text{"1_2.3."}, \text{"4_56.78."} \rangle$

outputs: $\langle \text{"1_2-3."}, \text{"4_56-78."} \rangle$

Example

-  Solution:

	Applicable?	Efficient?
Bottom-up Enumeration	O	X
Top-down Propagation	X	—
Duet	O	O

Existing Bottom-up Enumerative Strategy

- Enumerate expressions in order of increasing size
- Put smaller expressions together into larger ones

Size	Expressions
1	x “+” “_” “-” “.” 1 ... 9
2	Length(x) Length(“_”) Length(“-”) Length(“.”)
3	1 - 2 2 - 1 Concat(x , _) ... SubStr(x , 0, 1) ...
4	
...	

Existing Bottom-up Enumerative Strategy

Optimization: maintain only semantically unique expressions

Size	Expressions
1	x “+” “_” “-” “.” 1 ... 9
2	Length(x) Length(“_”) Length(“-”) Length(“.”)
3	1 - 2 2 - 1 ConCat(x , _) ... SubStr(x, 0, 1) ...
4	
...	

+ Generally applicable

- Limited scalability

Existing Bottom-up Enumerative Strategy

Optimization: maintain only semantically unique expressions

Size **Expressions**

The solution is too large to be quickly found by Bottom-up enumeration. ...

...

+ Generally applicable

- Limited scalability

Top-Down Propagation

- **Divide-and-conquer:** if a program $F(e_1, \dots, e_k)$ outputs O on some input, what should e_1, \dots, e_k output on the same input?

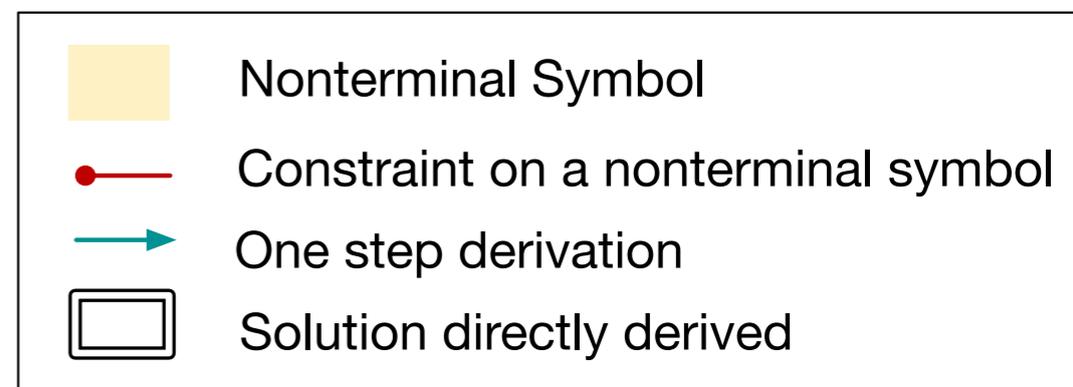
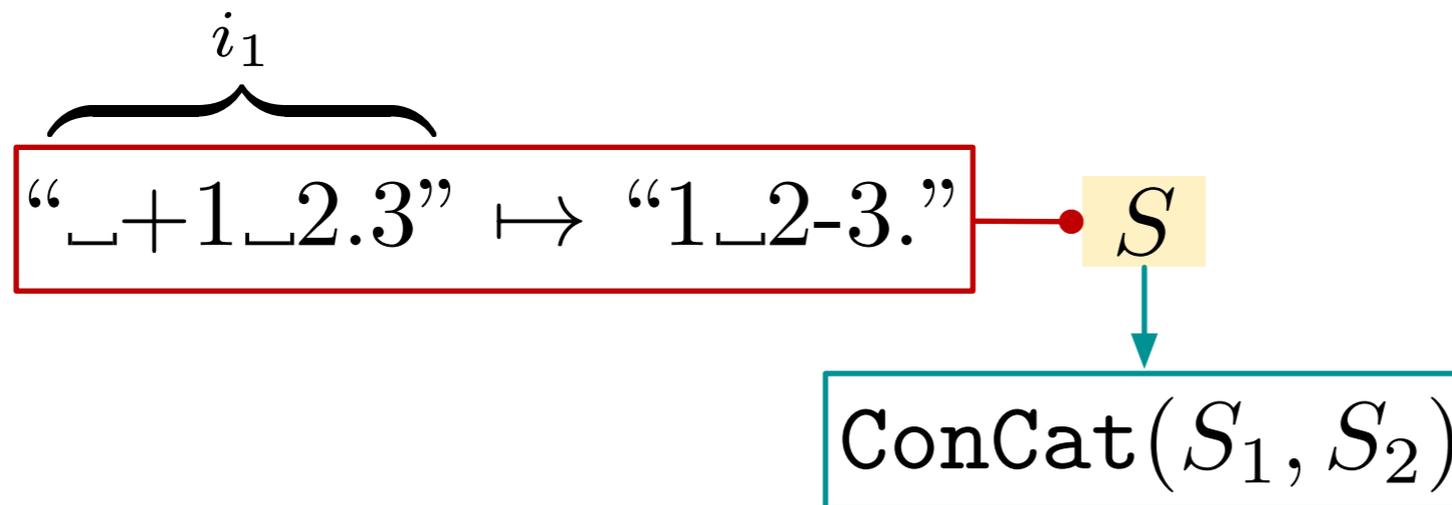
- infer specs for subexpressions using **inverse functions**

$$F^{-1}(o) = \{(a_1, \dots, a_k) \mid F(a_1, \dots, a_k) = o\}$$

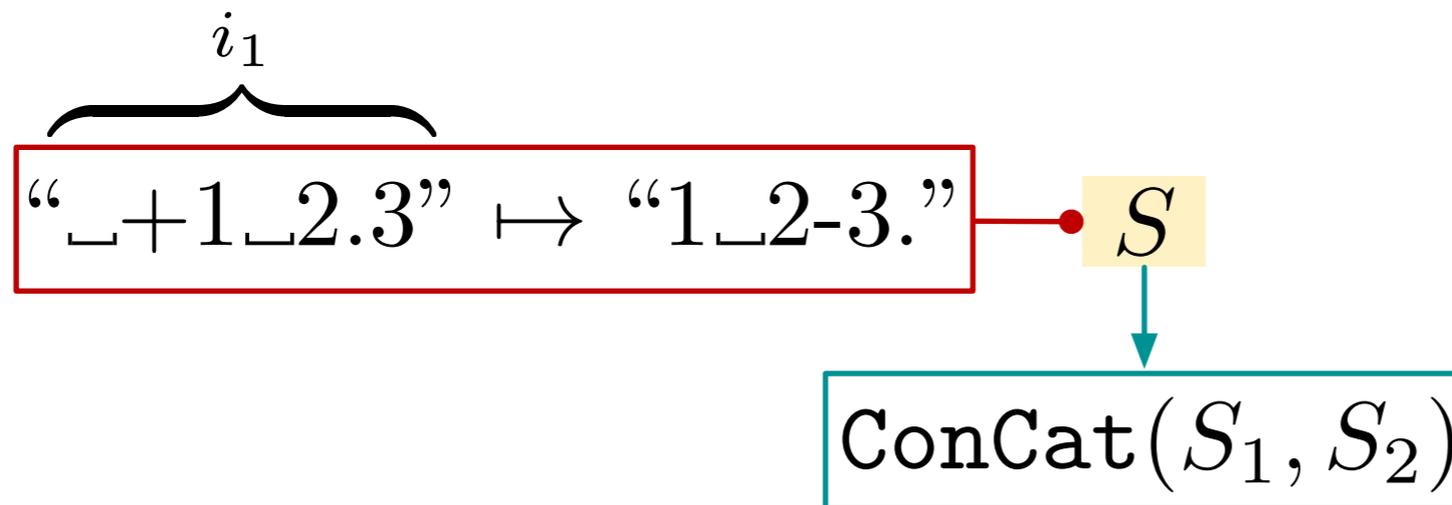
- *Inverse set (or pre-image):*

e.g., $\text{ConCat}^{-1}(\text{"USA"}) = \{(\text{"U"}, \text{"SA"}), (\text{"US"}, \text{"A"})\}$

Top-Down Propagation



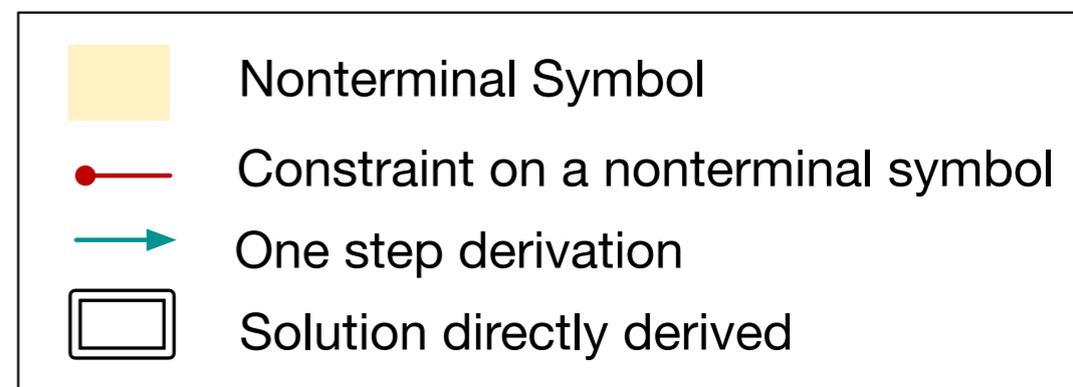
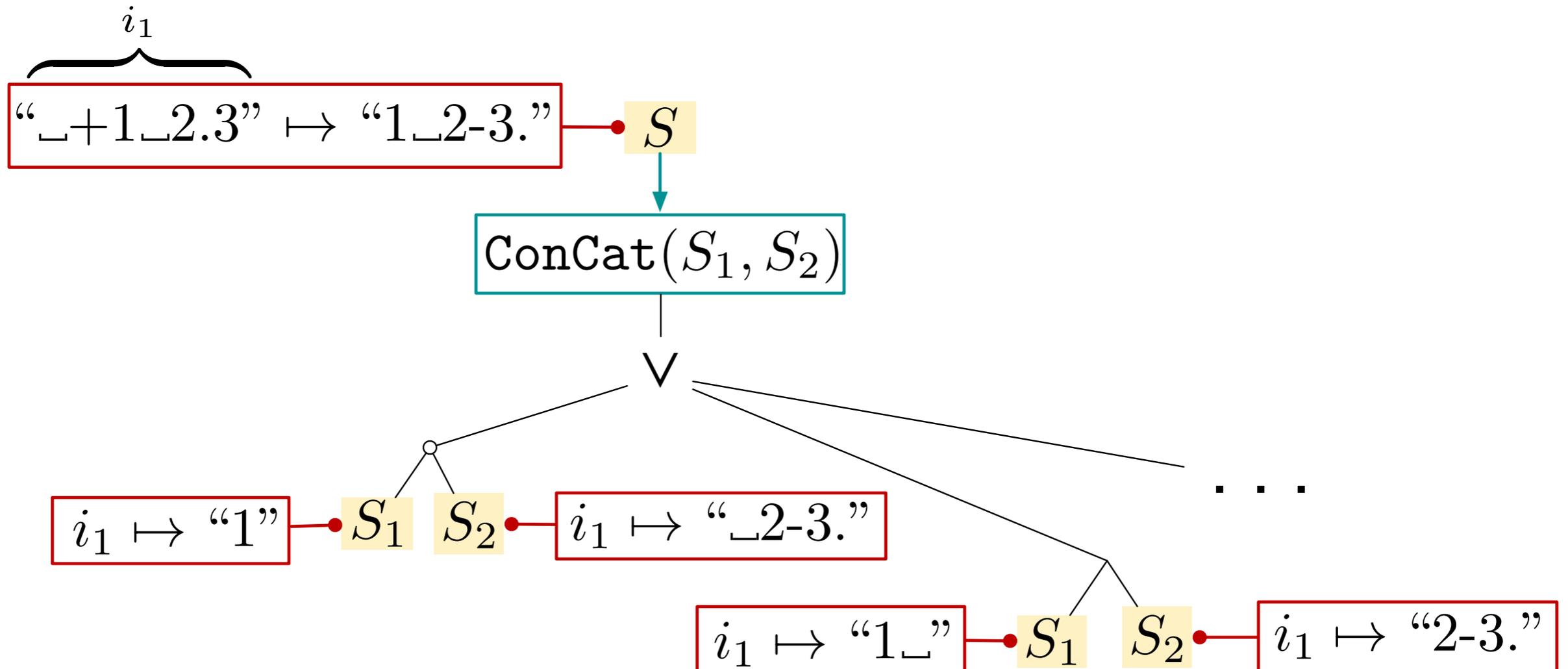
Top-Down Propagation



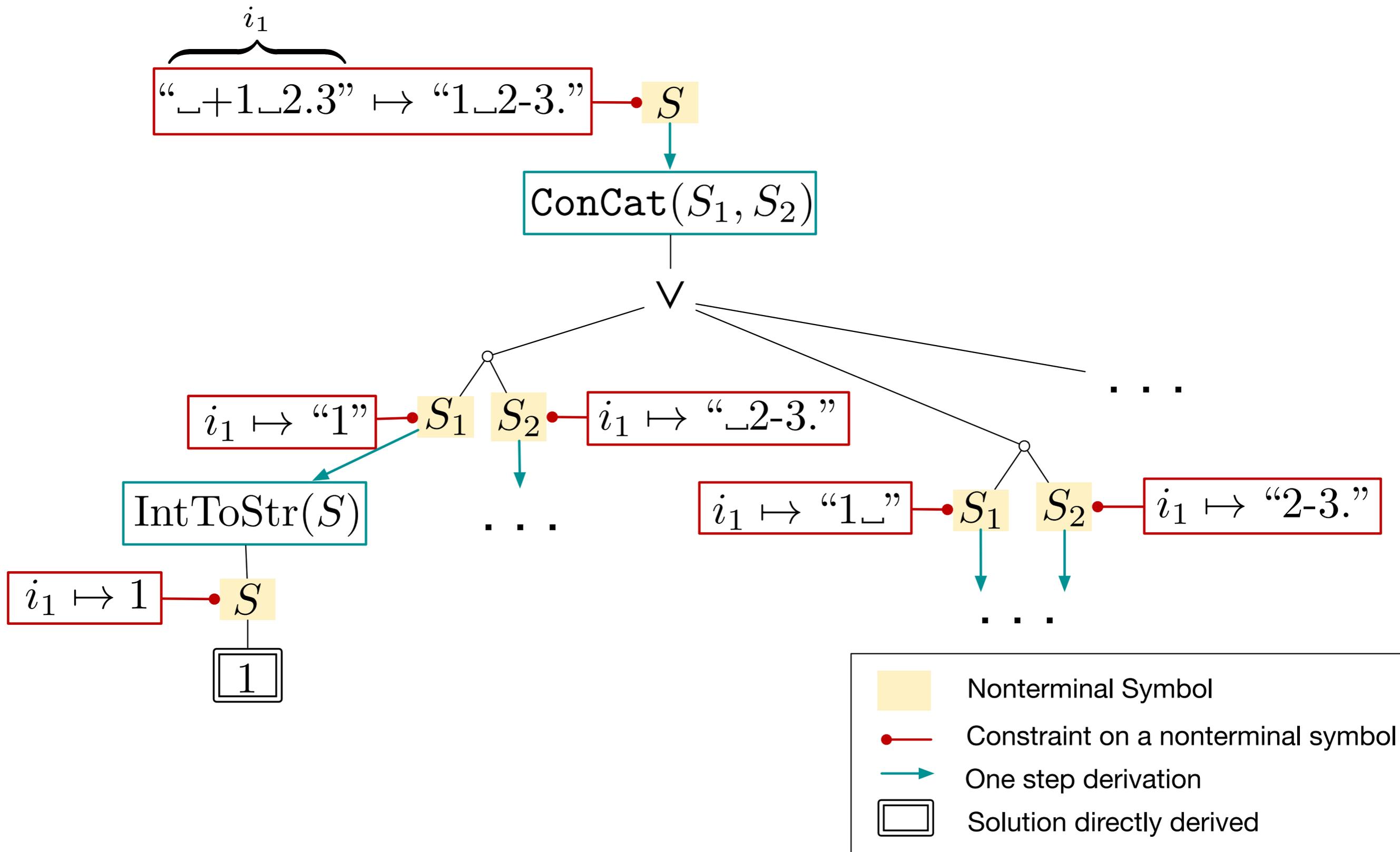
$$\text{ConCat}^{-1}(\text{"1␣2-3."}) = \{(\text{"1"}, \text{"␣2-3."}), (\text{"1␣"}, \text{"2-3."}), \dots\}$$

-  Nonterminal Symbol
-  Constraint on a nonterminal symbol
-  One step derivation
-  Solution directly derived

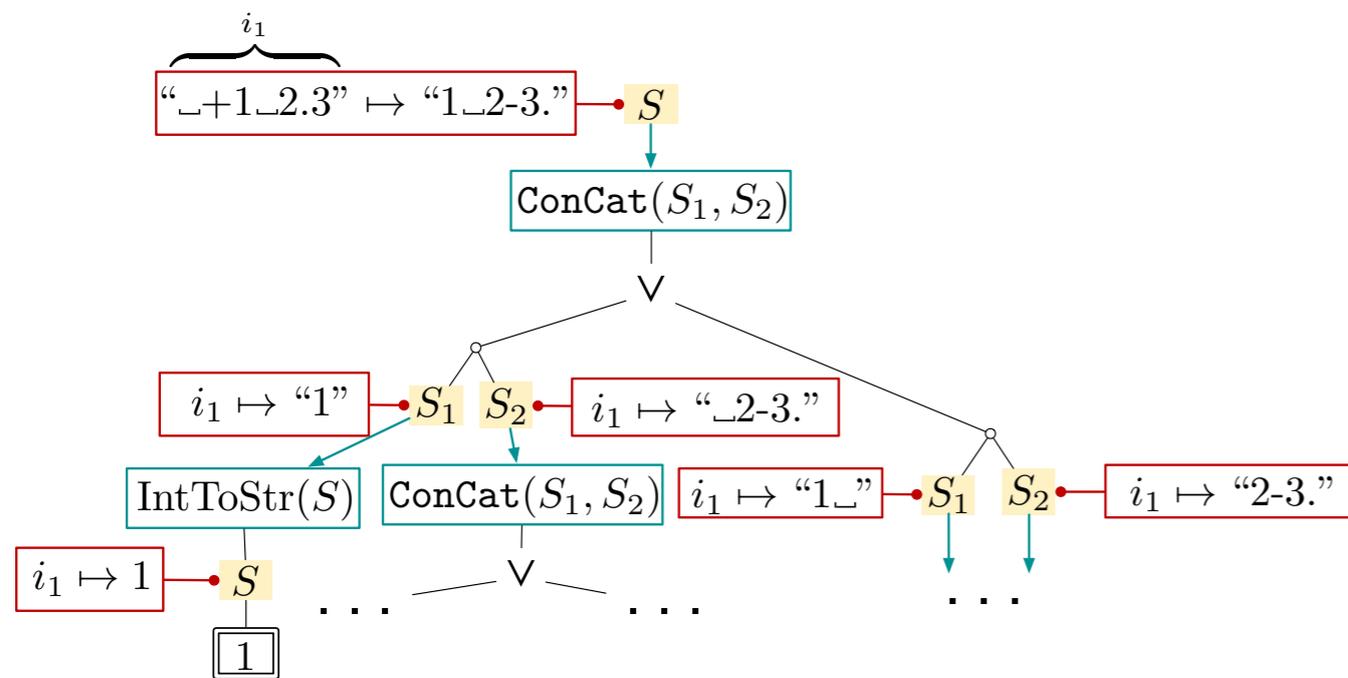
Top-Down Propagation



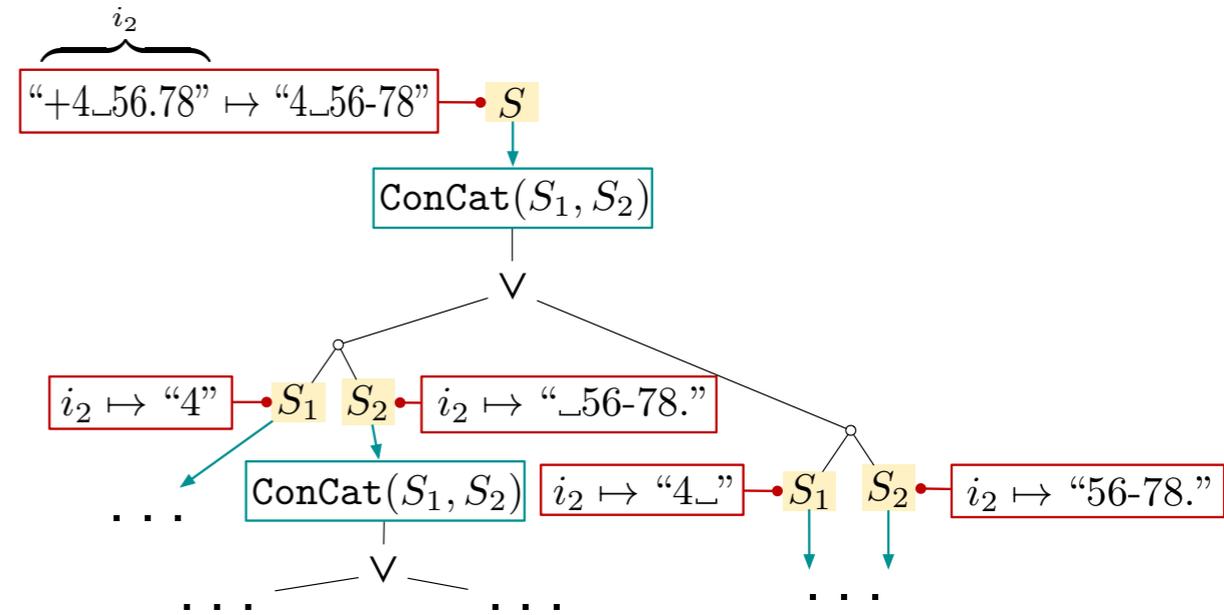
Top-Down Propagation



Top-Down Propagation



All possible solutions
for IO example 1



All possible solutions
for IO example 2

{Concat(Concat(IntToStr(1), “_”), ...),
Concat(Concat(IntToStr(1), Concat(“_” ...)),
...}



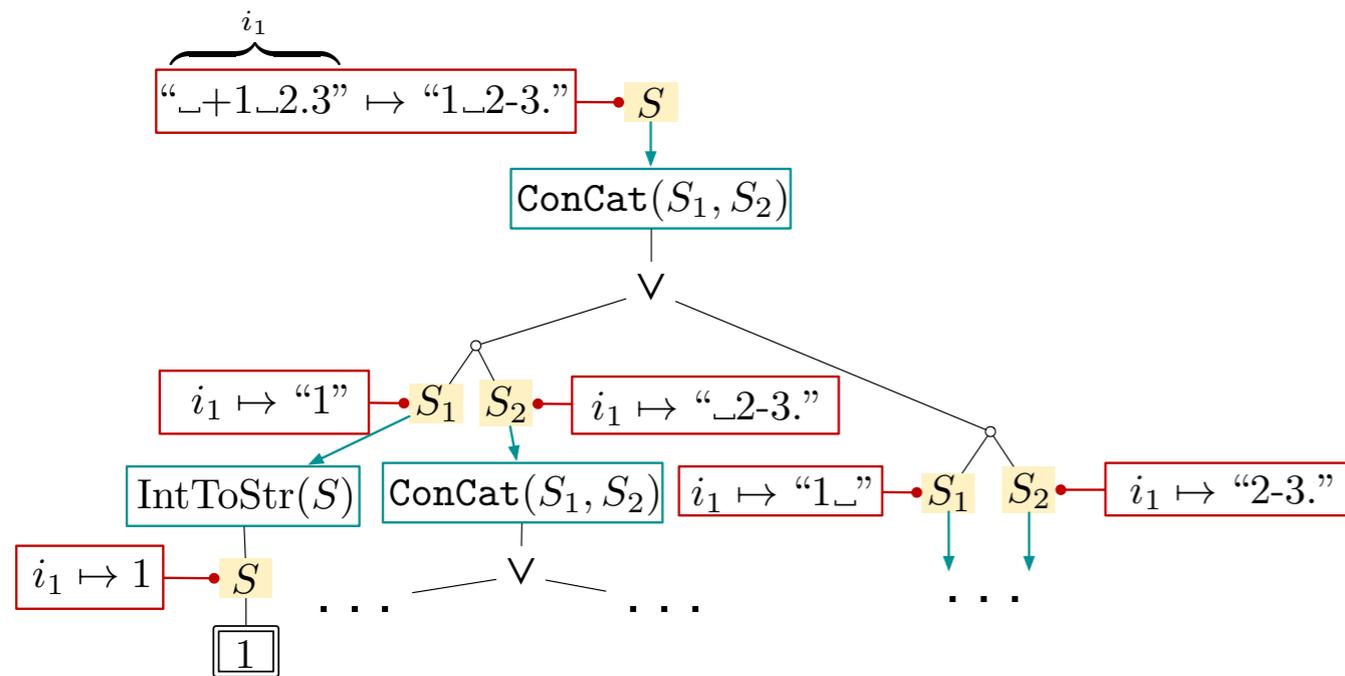
{Concat(Concat(IntToStr(4), “_”), ...),
Concat(Concat(IntToStr(4), Concat(“_” ...)),
...}

stored in a space-efficient data structure

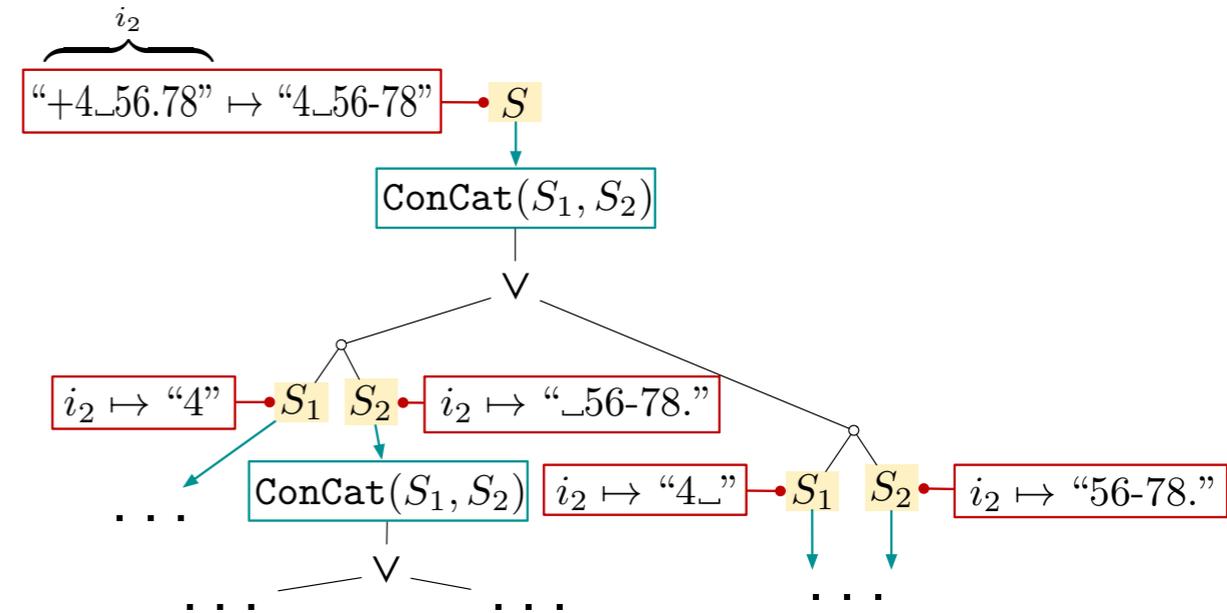


All possible final solutions

Top-Down Propagation



All possible solutions for IO example 1



All possible solutions for IO example 2

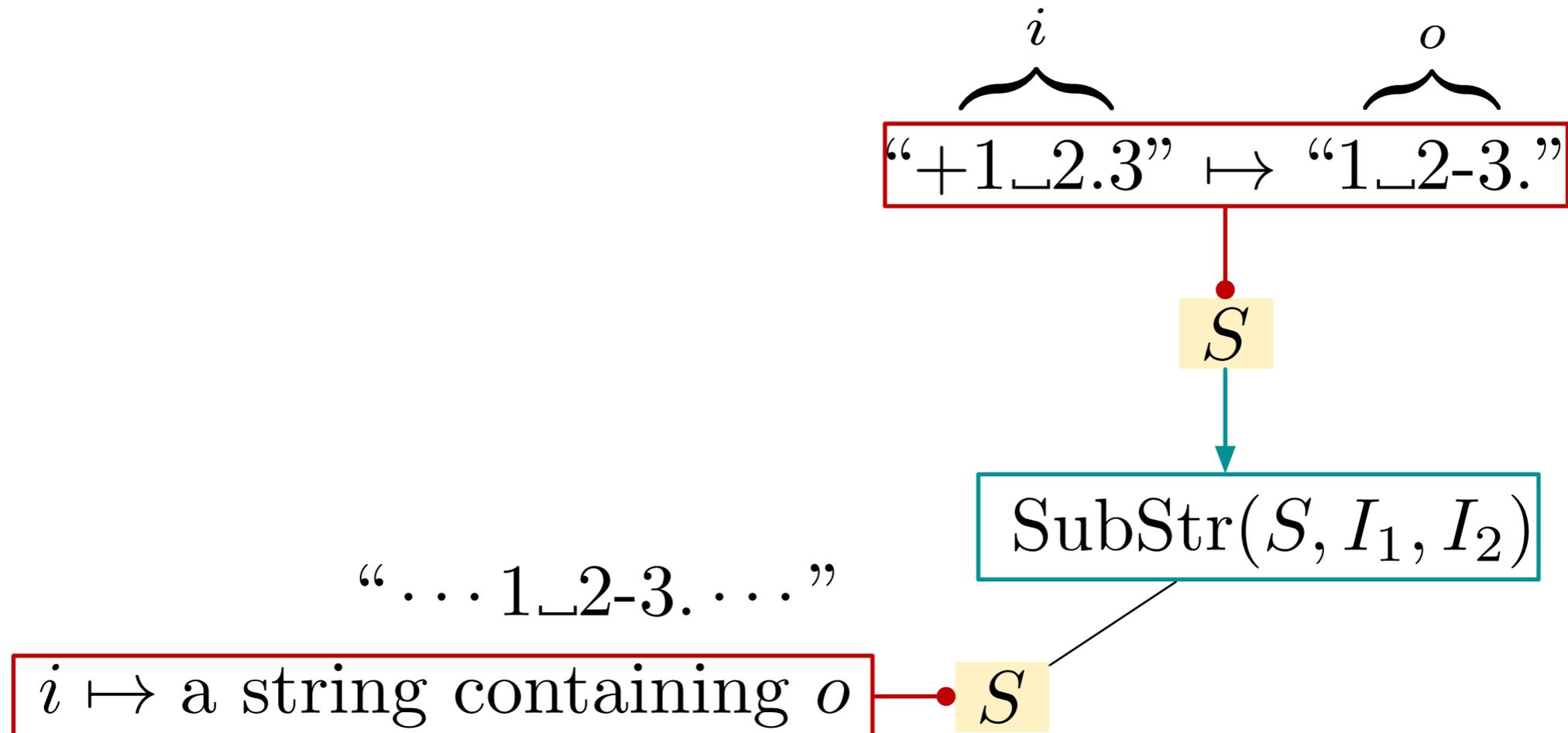
+ very efficient (divide-and-conquer, goal-directed)

- # of IO examples \uparrow \rightarrow performance \downarrow

- Not always applicable

Top-Down Propagation

Problem: ***unconstrained arguments***



Infinitely many such strings!
Inverse set SubStr^{-1} is infinite.

Top-Down Propagation

Problem: ***unconstrained arguments***

i o
“+1_2.3” \mapsto “1_2-3.”

Top-down propagation is not even applicable due to the general grammar.

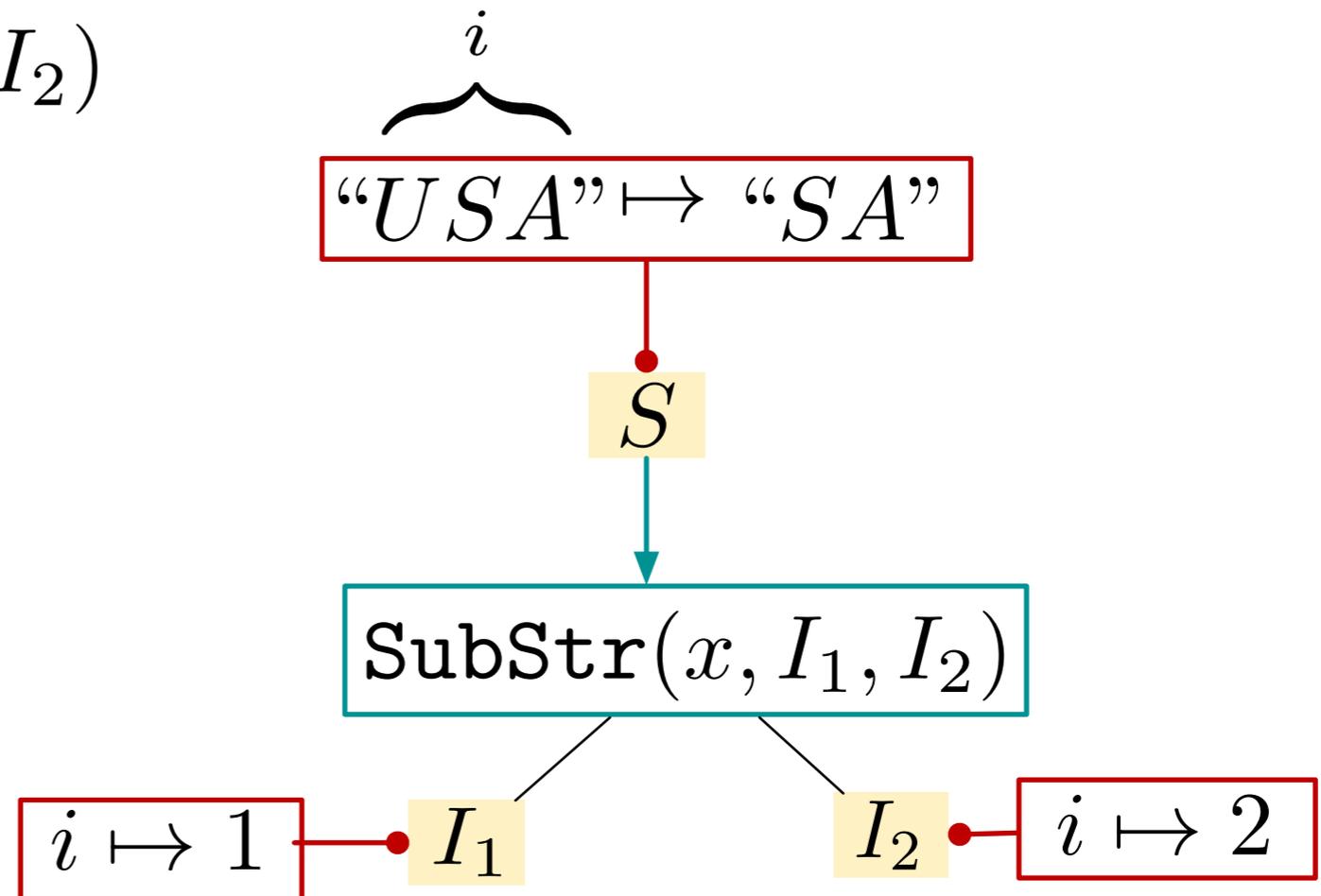
Infinitely many such strings!
Inverse set SubStr^{-1} is infinite.

Top-Down Propagation

Previous incomplete workaround: *restricting* the grammar (which limits the kinds of programs that can be synthesized)

$S \rightarrow \dots \mid \text{SubStr}(x, I_1, I_2)$

Input example



Top-Down Propagation

Flash Fill DSL

Tuple(String $x_1, \dots, \text{String } x_n$) \rightarrow String

top-level expr $T := C \mid \text{ifThenElse}(B, C, T)$

condition-free expr $C := A \mid \text{Concat}(A, C)$

atomic expression $A := \text{SubStr}(\boxed{X}, P, P) \mid \text{ConstantString}$

input string $X := x_1 \mid x_2 \mid \dots$

position expression $P := K \mid \text{Pos}(X, R_1, R_2, K)$

Only for DSLs with *Balanced Expressivity*

“DSL design = Art + *Lots* of iterations”

“The DSL should be *expressive enough* to represent various tasks... and *restricted enough* to allow efficient search.”

Gulwani et al., Program synthesis

Polozov et al., FlashMeta: a framework for inductive program synthesis

Gulwani et al., Programming by example (and its application to data wrangling)

Key Idea

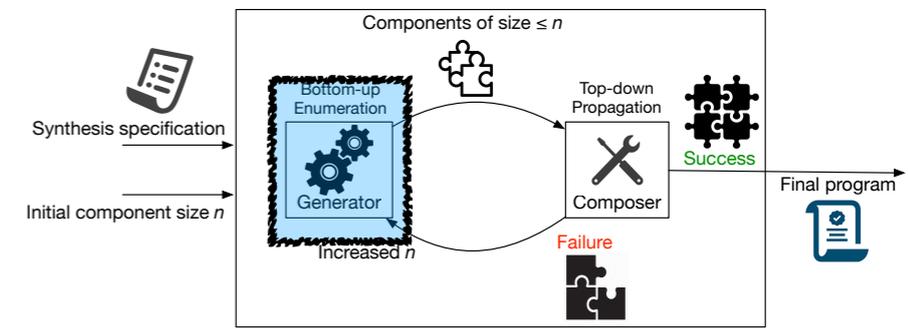
- Bottom-up enumeration can quickly identify *small* but *important* subexpressions of the solution.

Replace(SubStr(ConCat(x , "."), 2, Length(x) - 1), ".", "-")

In places of unconstrained arguments

- Our Top-down propagation: if an inverse set is infinite,
 - inverse set \subseteq {output of e | e is a component}, or
 - inverse set \subseteq {value similar to output of e | e is a component}

Component Generation

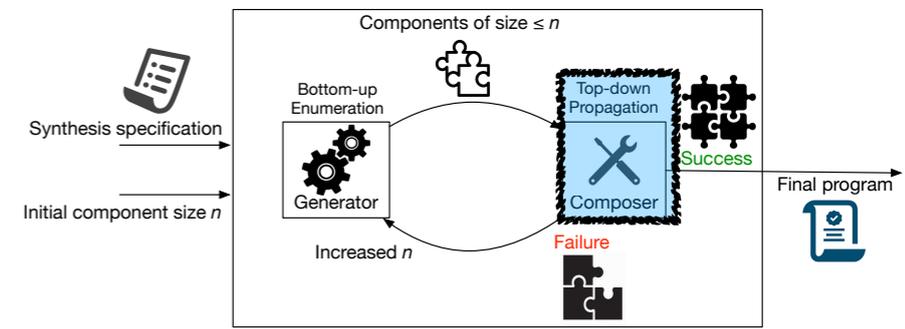


Suppose the initial component size is given **1**.

Bottom-up enumeration generates

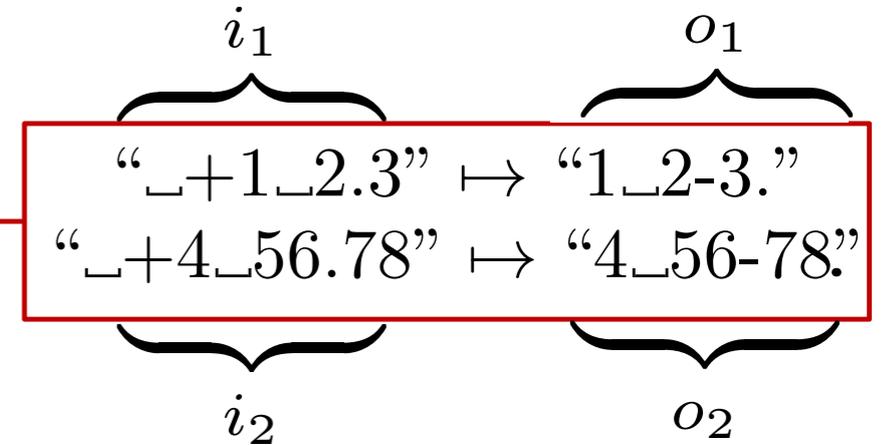
$$C = \left\{ \begin{array}{l} x, "+", "\cdot", "_", "_", \\ 1, \dots, 9 \end{array} \right\}$$

Inverse Set \subseteq Outputs of Components



Component Pool

$$C = \left\{ x, "+", ".", "_", "-", 1, \dots, 9 \right\}$$



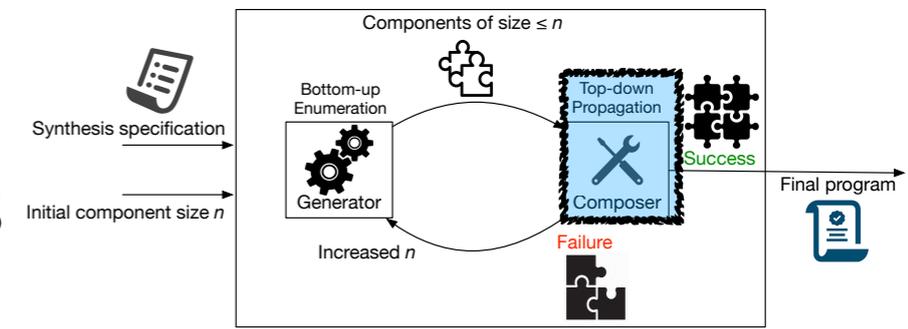
SubStr(S, I_1, I_2)

$$i_1 \mapsto \llbracket e \rrbracket(i_1) \quad i_2 \mapsto \llbracket e \rrbracket(i_2)$$

where $e \in C$, $o_1 \prec \llbracket e \rrbracket(i_1)$,
 $o_2 \prec \llbracket e \rrbracket(i_2)$

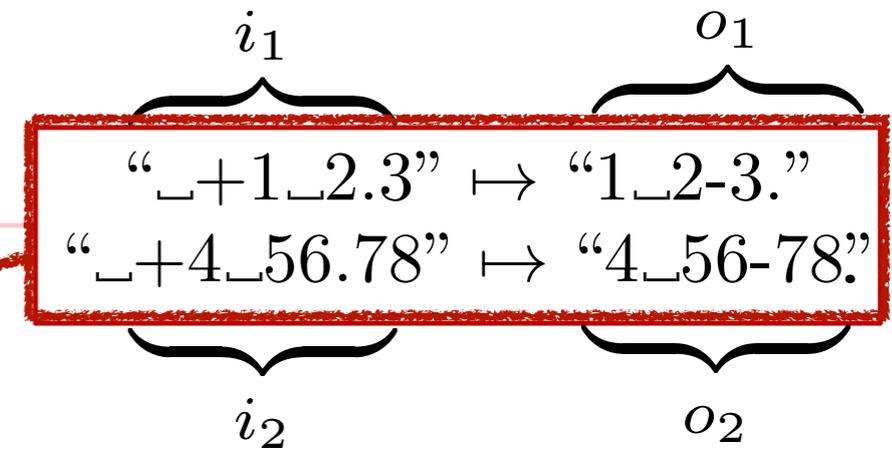
Search for a component that produces strings containing the outputs

Inverse Set \subseteq Outputs of Components



Component Pool

$$C = \left\{ x, "+", ".", "-", "_", \right. \\ \left. 1, \dots, 9 \right\}$$



SubStr(S, I_1, I_2)

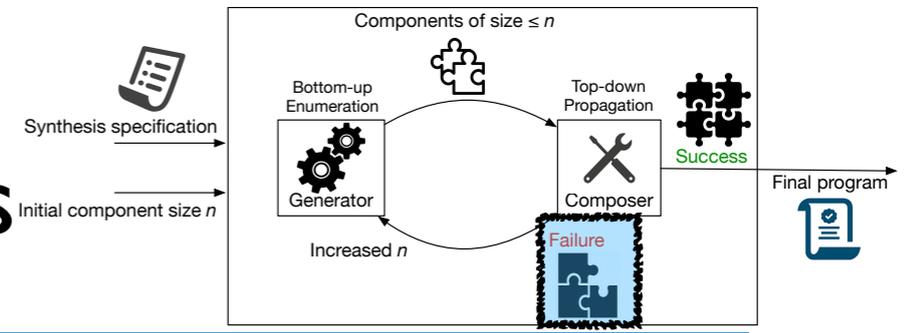
$$i_1 \mapsto \llbracket e \rrbracket(i_1) \quad i_2 \mapsto \llbracket e \rrbracket(i_2)$$

where $e \in C$, $o_1 \prec \llbracket e \rrbracket(i_1)$,
 $o_2 \prec \llbracket e \rrbracket(i_2)$

Simultaneous decomposition:

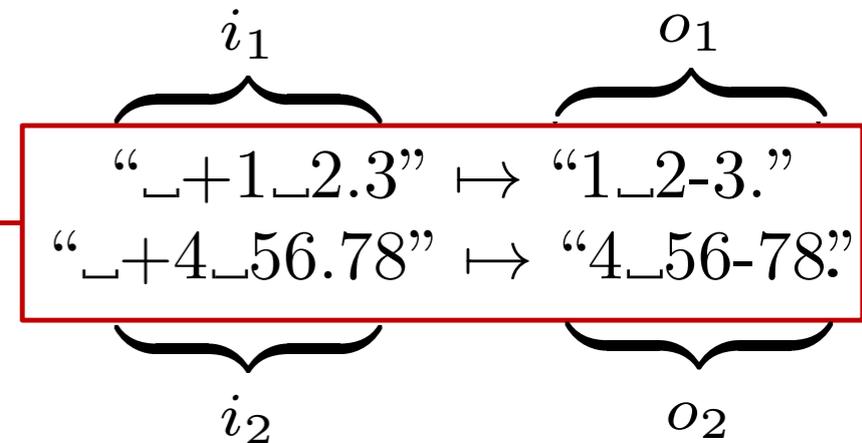
- handle multiple IO examples simultaneously
- no need for set intersection
- our method scales well on the # of IO examples.

Inverse Set \subseteq Outputs of Components



Component Pool

$$C = \left\{ x, "+", ".", "_", "-", 1, \dots, 9 \right\}$$



S

~~SubStr(S, I_1, I_2)~~

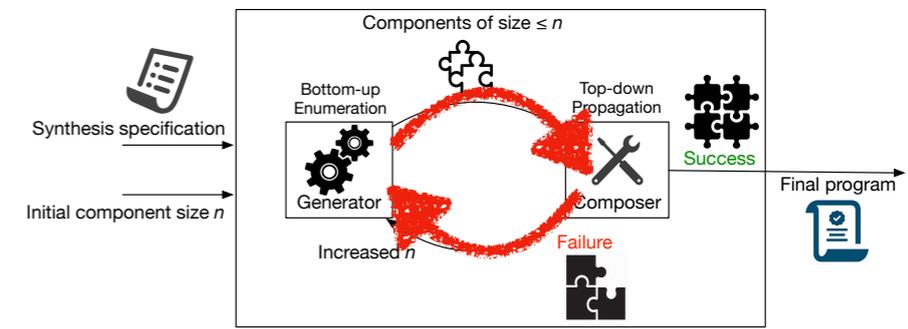
S

$i_1 \mapsto \llbracket e \rrbracket(i_1)$ $i_2 \mapsto \llbracket e \rrbracket(i_2)$
 where $e \in C$, $o_1 \prec \llbracket e \rrbracket(i_1)$,
 $o_2 \prec \llbracket e \rrbracket(i_2)$

$\nexists e \in C. o_1 \prec \llbracket e \rrbracket(i_1) \wedge o_2 \prec \llbracket e \rrbracket(i_2)$

**No such a component.
Reject the hypothesis.**

We don't miss a solution



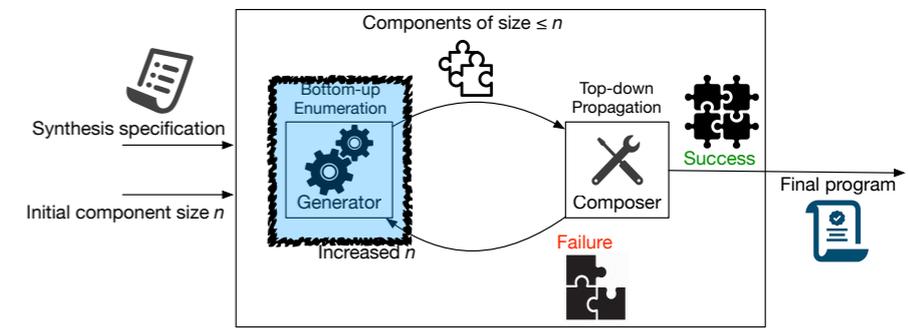
- We may miss the following solution of form $\text{SubStr}(\dots)$ at this iteration because we used size 1 components.

$\text{SubStr}(\text{Replace}(\text{Concat}(x, "."), ".", "-"), 2, \text{Length}(x) - 1)$

We should've used this component (size 6)

- But we will find it by increasing the component size.
- **Search completeness:** if a solution exists, we find it.

Component Generation

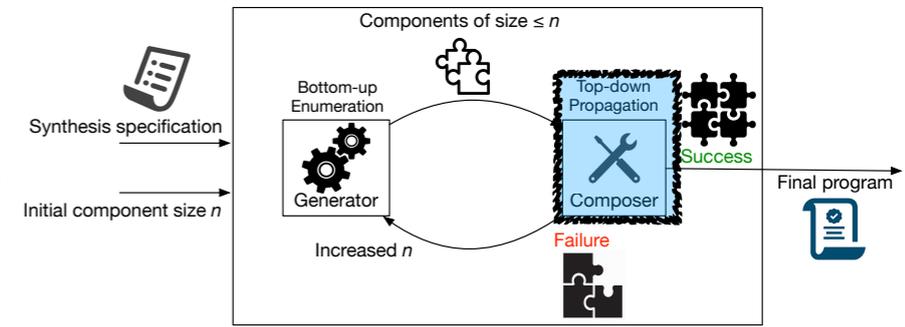


Now suppose the component size is increased to 3.

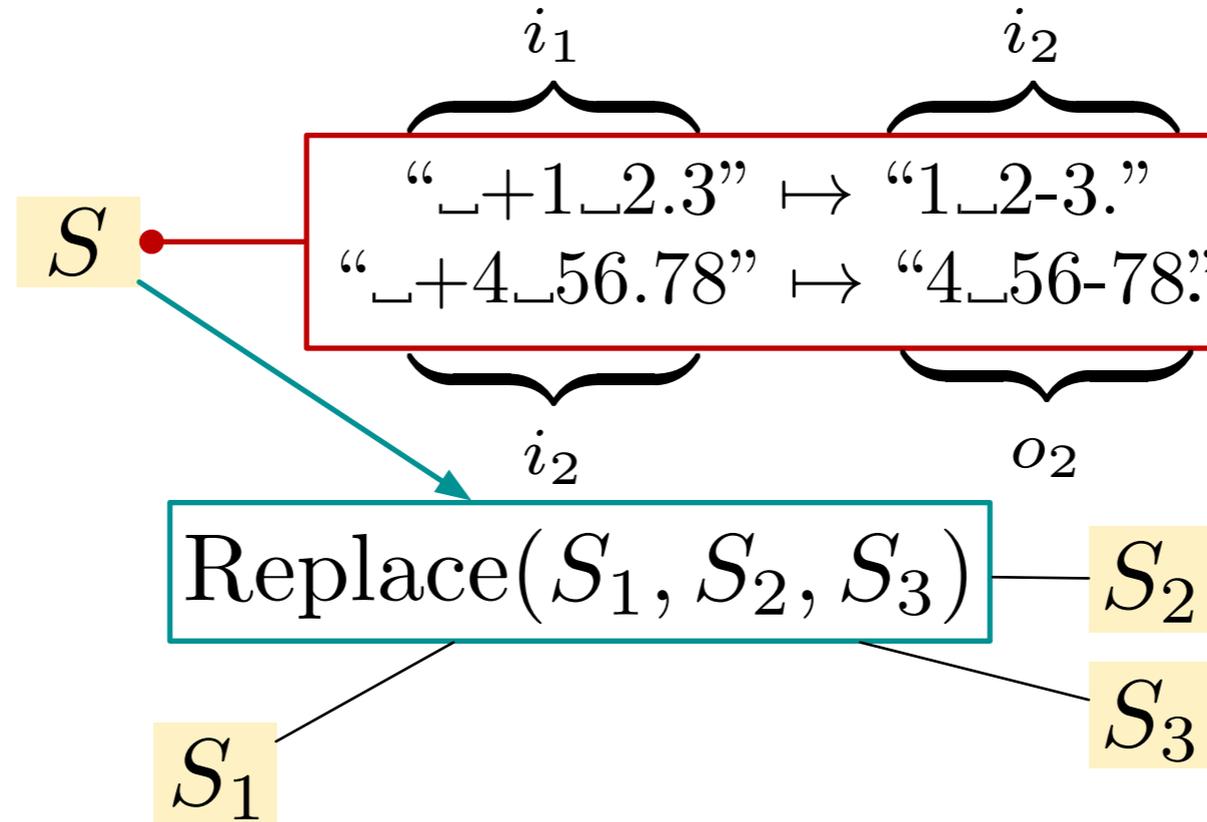
Bottom-up enumeration generates

$$C = \left\{ \begin{array}{l} x, \text{“.”}, \text{“_”}, \text{“-”}, \\ 1, \dots, \text{Length}(x), \\ \text{ConCat}(x, \text{“.”}), \\ \dots \end{array} \right\}$$

Problem of "Inverse Set \subseteq Outputs of Components"



- Suppose now we consider the hypothesis $\text{Replace}(S_1, S_2, S_3)$



- Simple Replace^{-1} similar to SubStr^{-1} needs the following component

$\text{Replace}(\text{SubStr}(\text{ConCat}(x, "."), 2, \text{Length}(x) - 1), ".", "-")$

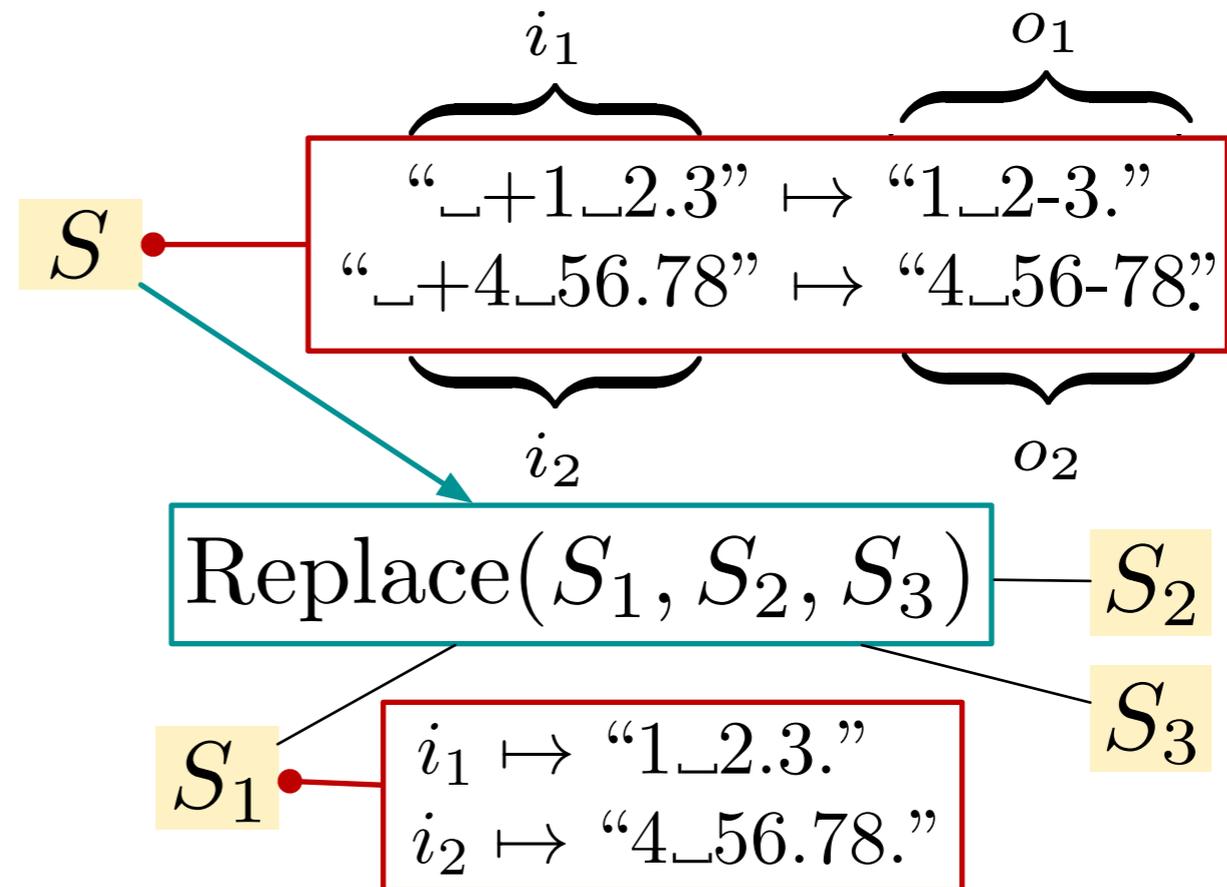
which is large (size 9) and cannot be efficiently generated.

Component-guided Inverse Function

- We generate subproblems that can be eventually solved by the current small components.
- Our Replace^{-1} : find a component producing strings most similar to the desired outputs $\rightarrow \text{ConCat}(x, \text{"."})$
- Compute *alignments* and generate specs for arguments

$\llbracket \text{ConCat}(x, \text{"."}) \rrbracket(i_1)$	␣	+	1	␣	2	.	3	.		
o_1 (desired output 1)	€	€	1	␣	2	-	3	.		
$\llbracket \text{ConCat}(x, \text{"."}) \rrbracket(i_2)$	␣	+	4	␣	5	6	.	7	8	.
o_2 (desired output 2)	€	€	4	␣	5	6	-	7	8	.

Component-guided Inverse Function

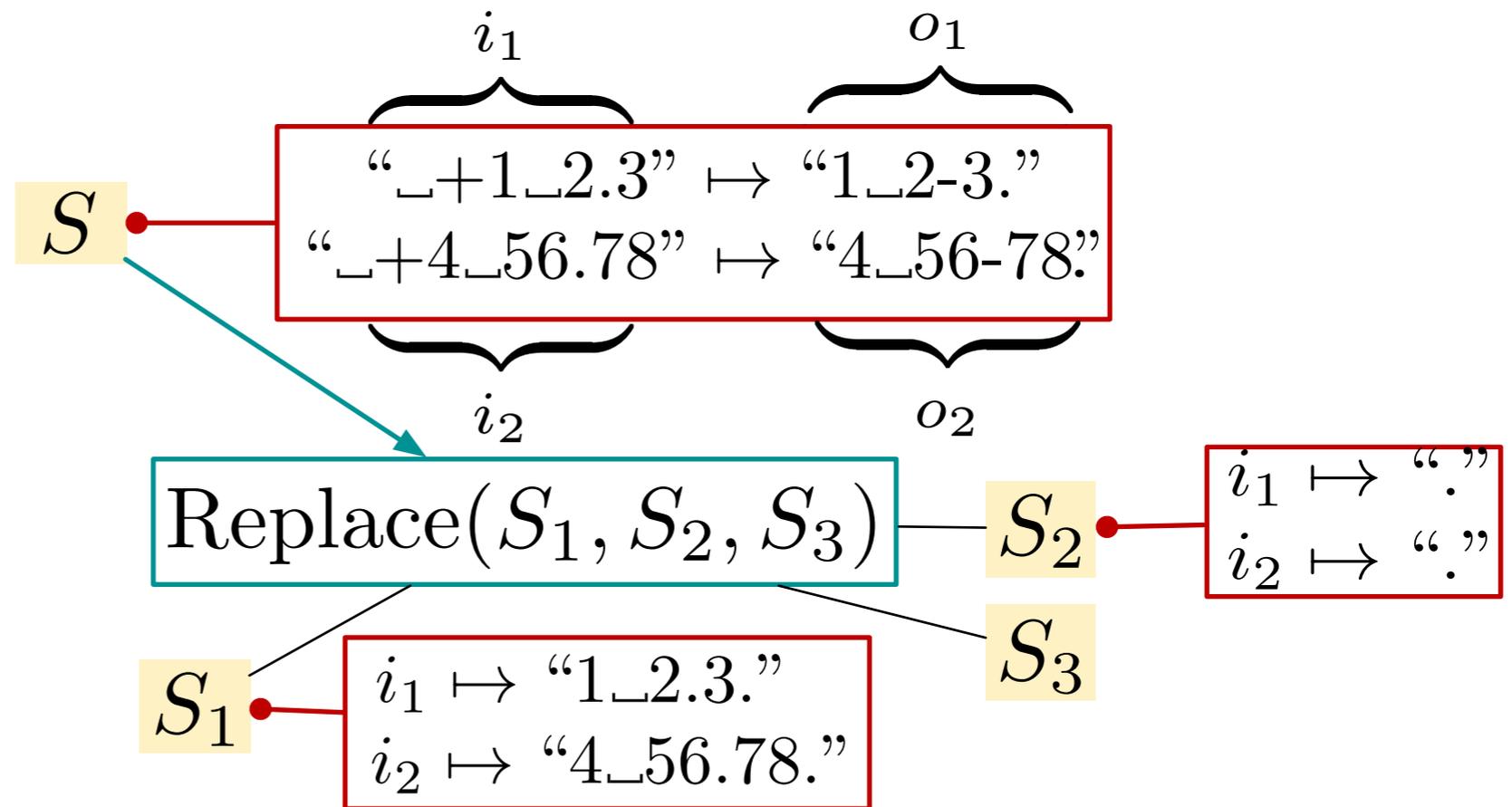


_	+	1	_	2	.	3	.
€	€	1	_	2	-	3	.

Desired outputs for S_1

_	+	4	_	5	6	.	7	8	.
€	€	4	_	5	6	-	7	8	.

Component-guided Inverse Function

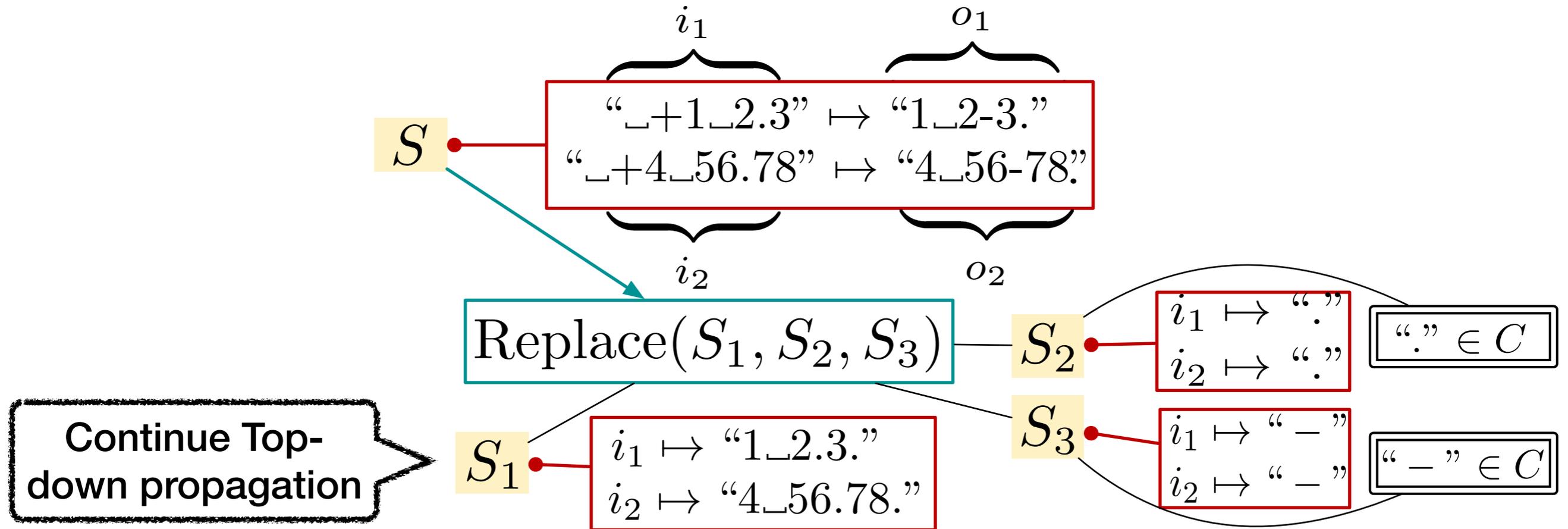


␣	+	1	␣	2	.	3	.
€	€	1	␣	2	-	3	.

Desired output for S_2

␣	+	4	␣	5	6	.	7	8	.
€	€	4	␣	5	6	-	7	8	.

Component-guided Inverse Function



$_$	$+$	1	$_$	2	$.$	3	$.$
ϵ	ϵ	1	$_$	2	$-$	3	$.$

Desired output for S_3

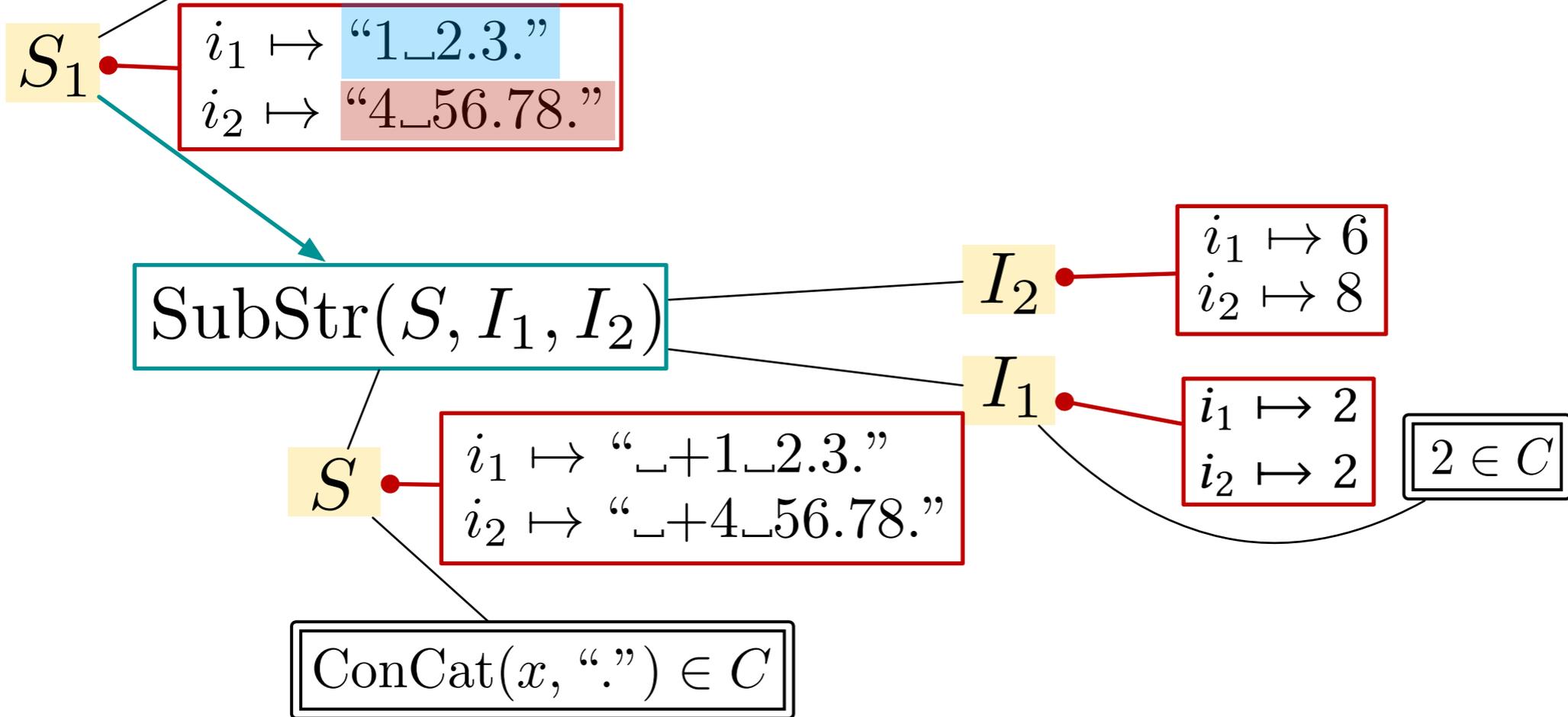
$_$	$+$	4	$_$	5	6	$.$	7	8	$.$
ϵ	ϵ	4	$_$	5	6	$-$	7	8	$.$

Component-guided Inverse Function

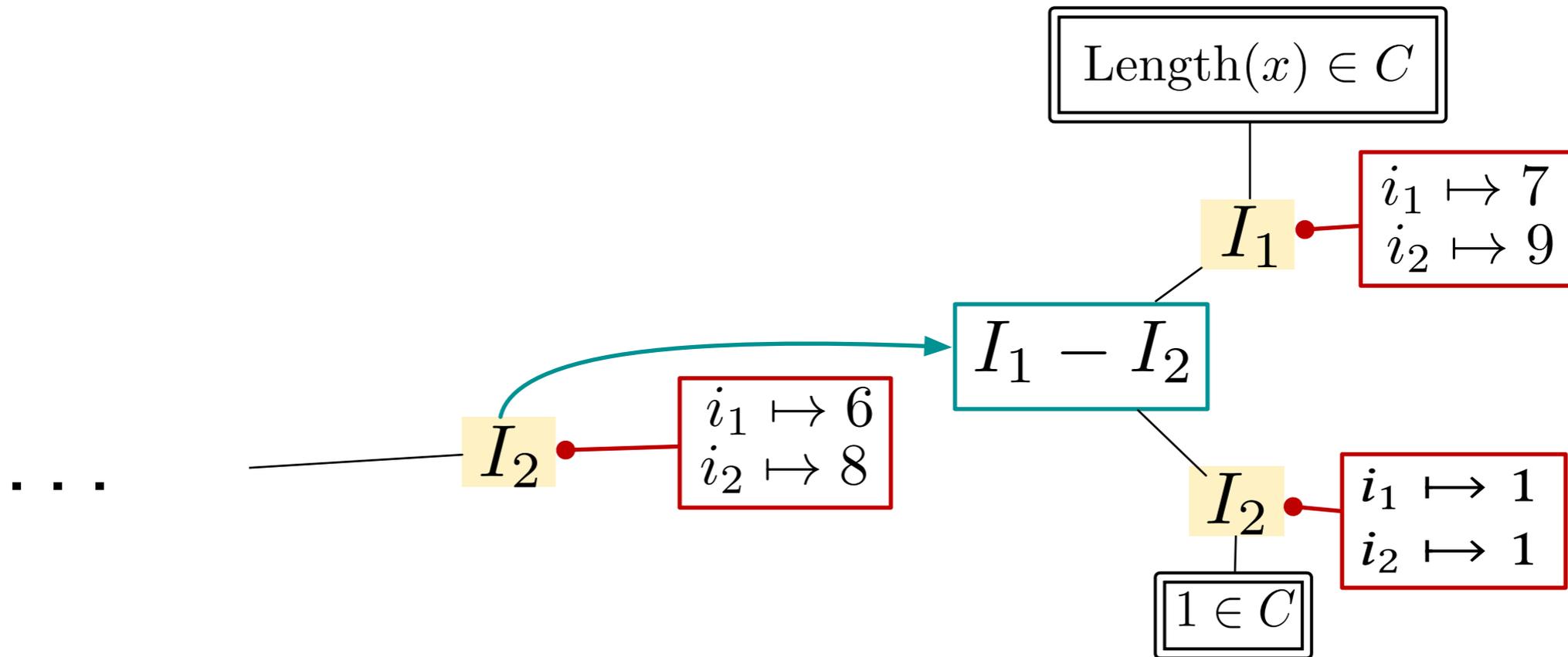
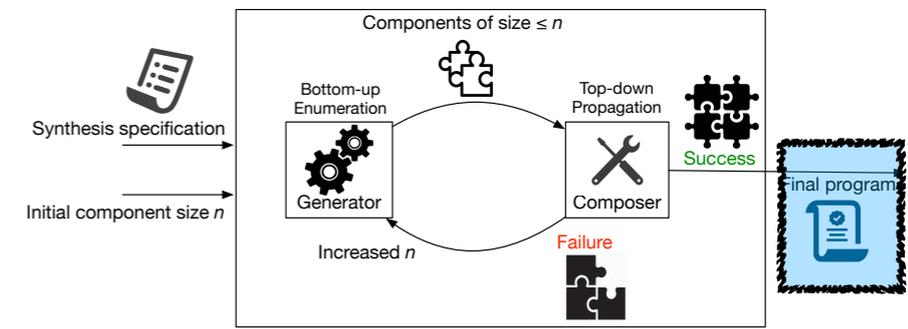
$\llbracket \text{ConCat}(x, ".") \rrbracket(i_1)$ $_$ | $+$ | 1 | $_$ | 2 | $.$ | 3 | $.$

...

$\llbracket \text{ConCat}(x, ".") \rrbracket(i_2)$ $_$ | $+$ | 4 | $_$ | 5 | 6 | $.$ | 7 | 8 | $.$



Component-guided Inverse Function



No subproblems unsolved. The solution is found.

Our Inverse Functions for SyGuS

- Specialized for the operators supported in the SyGuS specification language (strings, bit vectors, LIA, and SAT)
- Type 1: deducing specs **not directly solvable** with components (e.g., Replace⁻¹) → consecutive propagations
- Type 2: deducing specs **directly solvable** with components (e.g., SubStr⁻¹) → solved in a single step

Universal Inverse Function

- In cases where (1) no particular domain knowledge of an underlying operator is exploitable, or (2) computing pre-images is expensive, we use the *universal inverse function*.
- Suppose a spec $i \mapsto o$ is given on N and $N \rightarrow F(N_1, \dots, N_k)$
- Then the universal inverse function computes the following pre-images:

enumerate all possible combinations of components

$$F^{-1}(o) = \{(\llbracket e_1 \rrbracket(i), \dots, \llbracket e_k \rrbracket(i)) \mid e_1, \dots, e_k \in C, \llbracket F(e_1, \dots, e_k) \rrbracket(i) = o\}$$

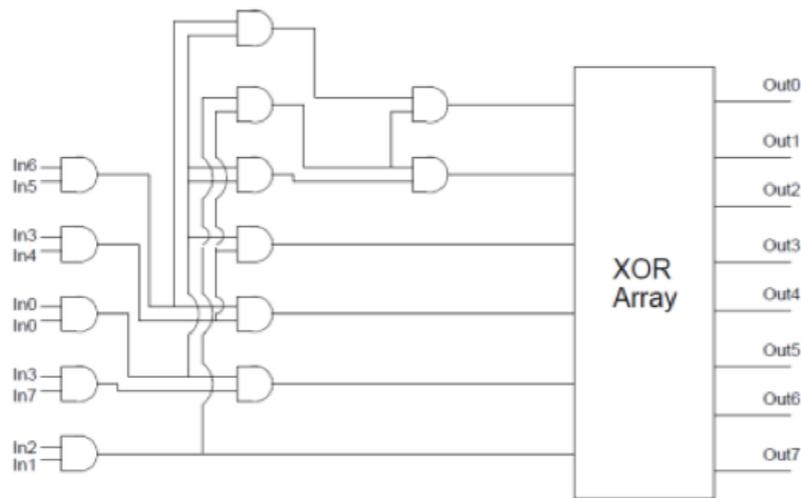
Evaluation Setup

- Benchmarks: **1,536** SyGuS problems
 - **1,167** from the SyGuS annual competitions + **369** from optimization tasks for homomorphic evaluation [Lee et al. PLDI'20]
- Comparison to three baselines (Timeout 1 hour):
 - EUSolver: winner of 2016 SyGuS competition
 - CVC4: winner of 2017 - 2019 SyGuS competition
 - Euphony [Lee et al. PLDI'18]: statistical model-guided synthesizer

Benchmarks

	A	B
1	Email	Column 2
2	Nancy.FreeHafer@fourthcoffee.com	nancy freehafer
3	Andrew.Cencici@northwindtraders.com	andrew cencici
4	Jan.Kotas@litwareinc.com	jan kotas
5	Mariya.Sergienko@gradicdesigninstitute.com	mariya sergienko
6	Steven.Thorpe@northwindtraders.com	steven thorpe
7	Michael.Neipper@northwindtraders.com	michael neipper
8	Robert.Zare@northwindtraders.com	robert zare
9	Laura.Giussani@adventure-works.com	laura giussani
10	Anne.HL@northwindtraders.com	anne hl
11	Alexander.David@contoso.com	alexander david
12	Kim.Shane@northwindtraders.com	kim shane
13	Manish.Chopra@northwindtraders.com	manish chopra
14	Gerwald.Oberleitner@northwindtraders.com	gerwald oberleitner
15	Amr.Zaki@northwindtraders.com	amr zaki
16	Yvonne.McKay@northwindtraders.com	yvonne mckay
17	Amanda.Pinto@northwindtraders.com	amanda pinto

STRING: End-user programming
205 problems



CIRCUIT: Attack-resilient crypto circuits + Optimized homomorphic evaluation circuits
581 problems

complement

```
~ 01010001110101110000000000001111
   1010111000101000111111111110000
```

bitwise and

```
01010001110101110000000000001111
& 00110001011011100011000101101110
   0001000101000110000000000001110
```

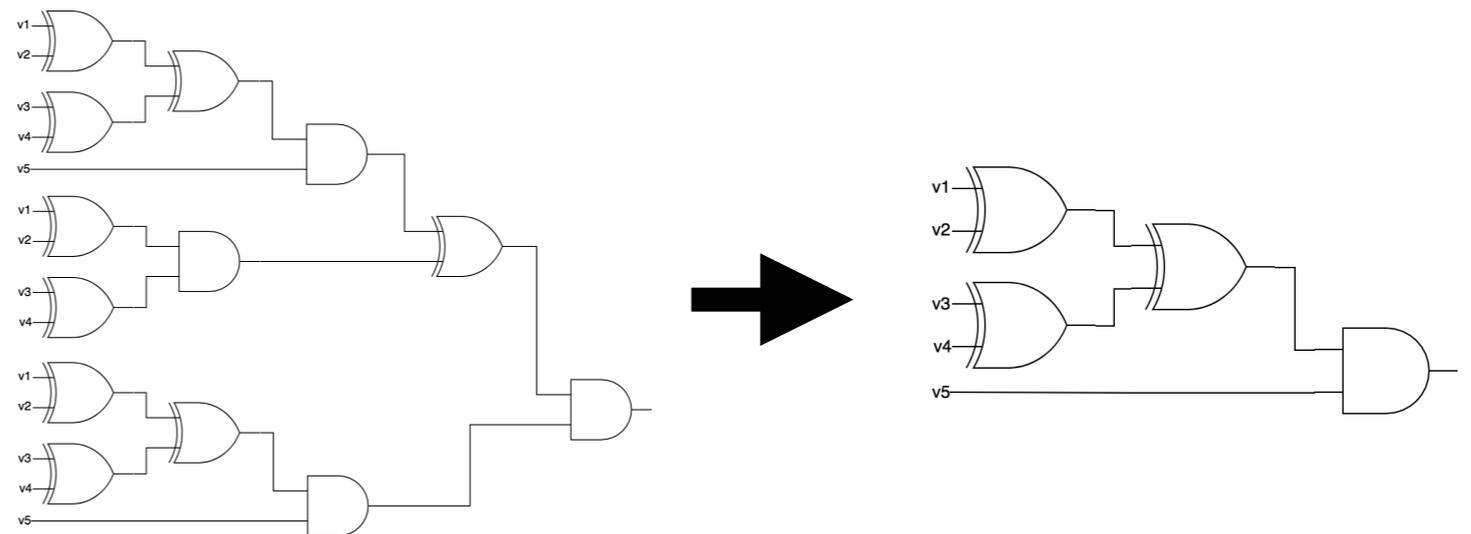
bitwise or

```
01010001110101110000000000001111
| 00110001011011100011000101101110
   011100011111111110011000101101111
```

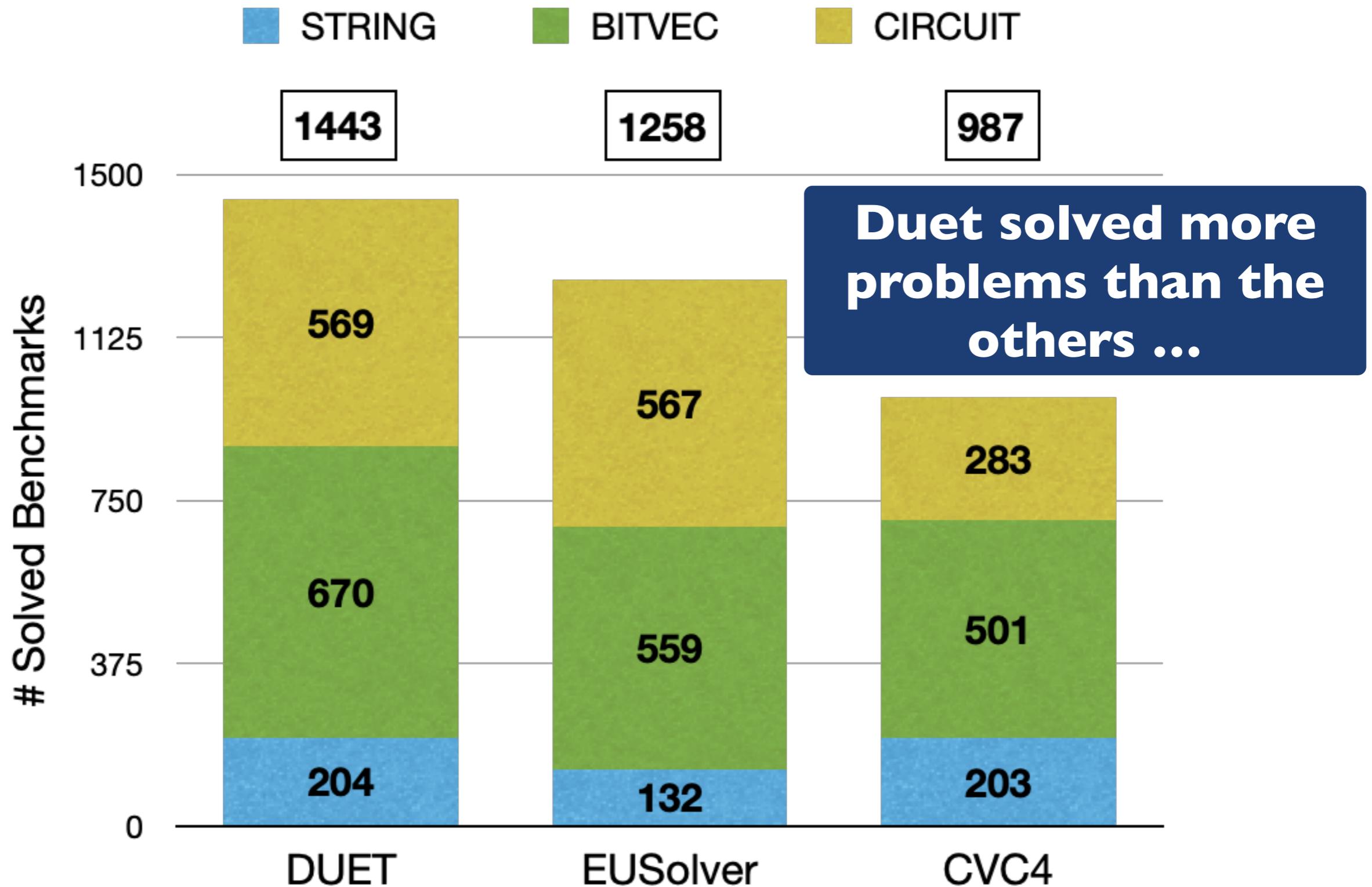
bitwise xor

```
01010001110101110000000000001111
^ 00110001011011100011000101101110
   01100000101110010011000101100001
```

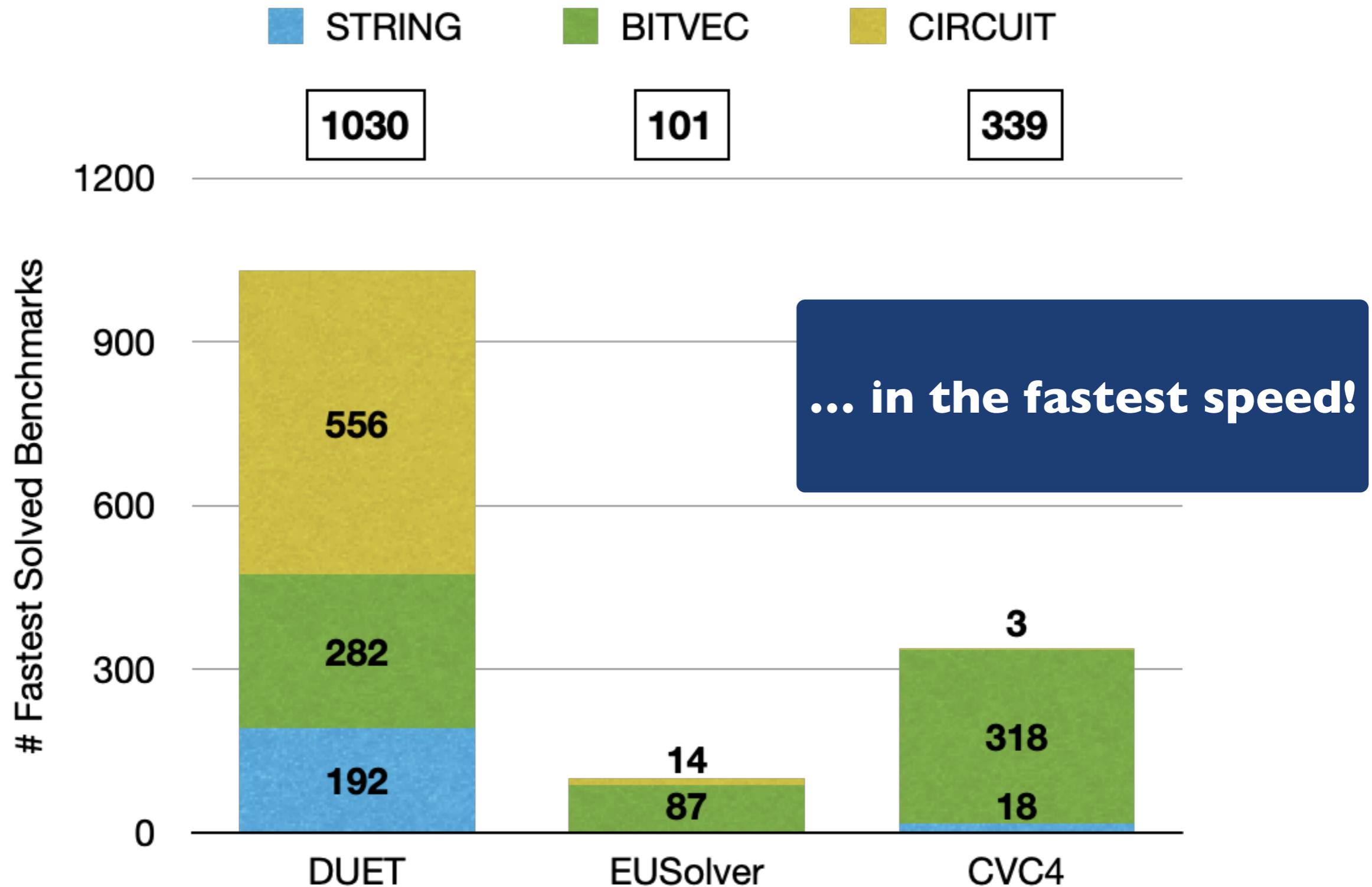
BITVEC: Efficient low-level algorithm
750 problems



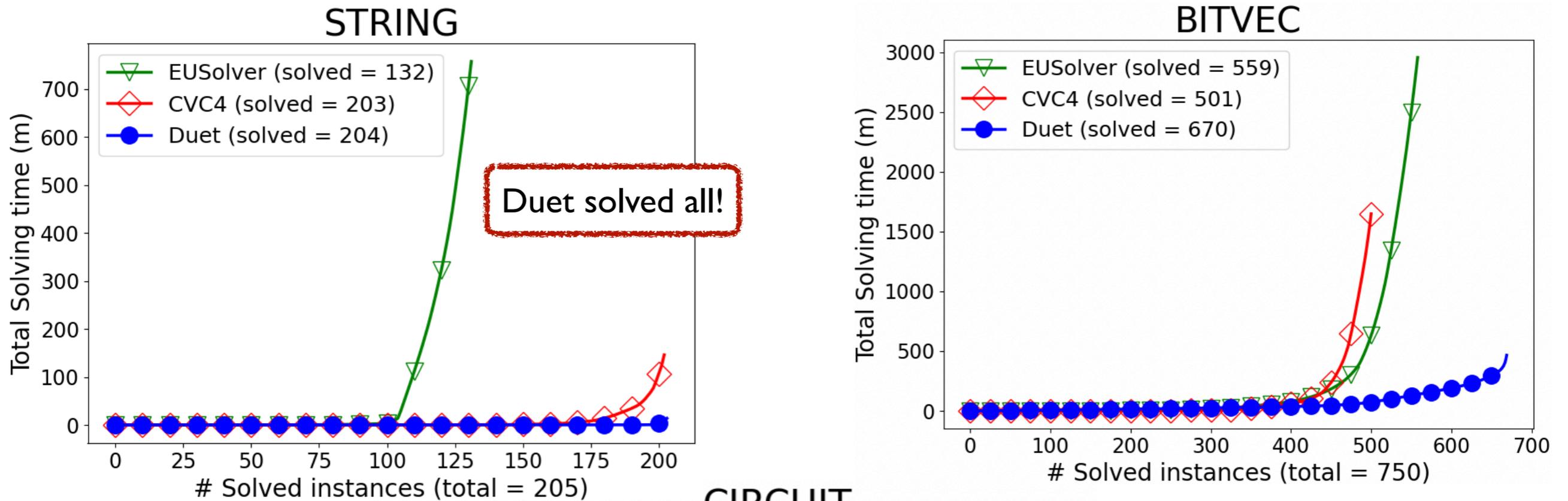
Comparison to the Winners of SyGuS Competitions



Comparison to the Winners of SyGuS Competitions

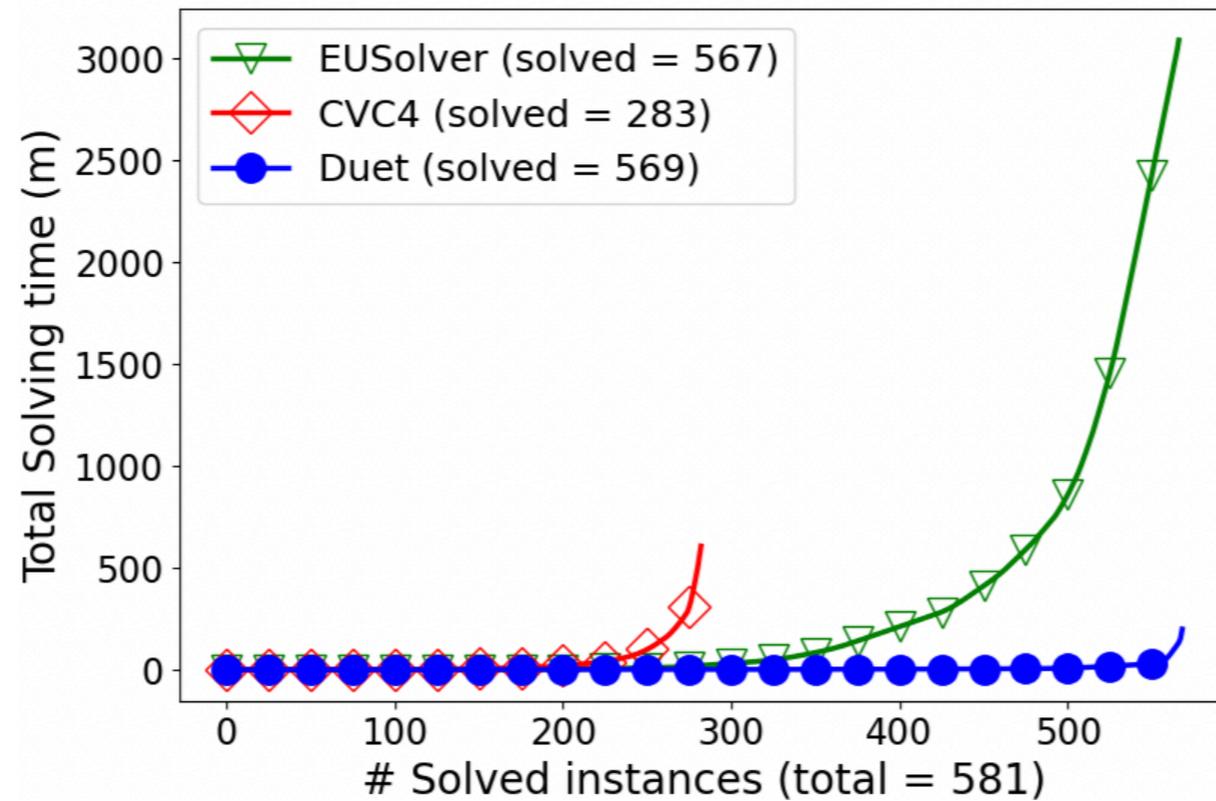


Comparison to the Winners of SyGuS Competitions



Duet solved all!

CIRCUIT



Duet solved

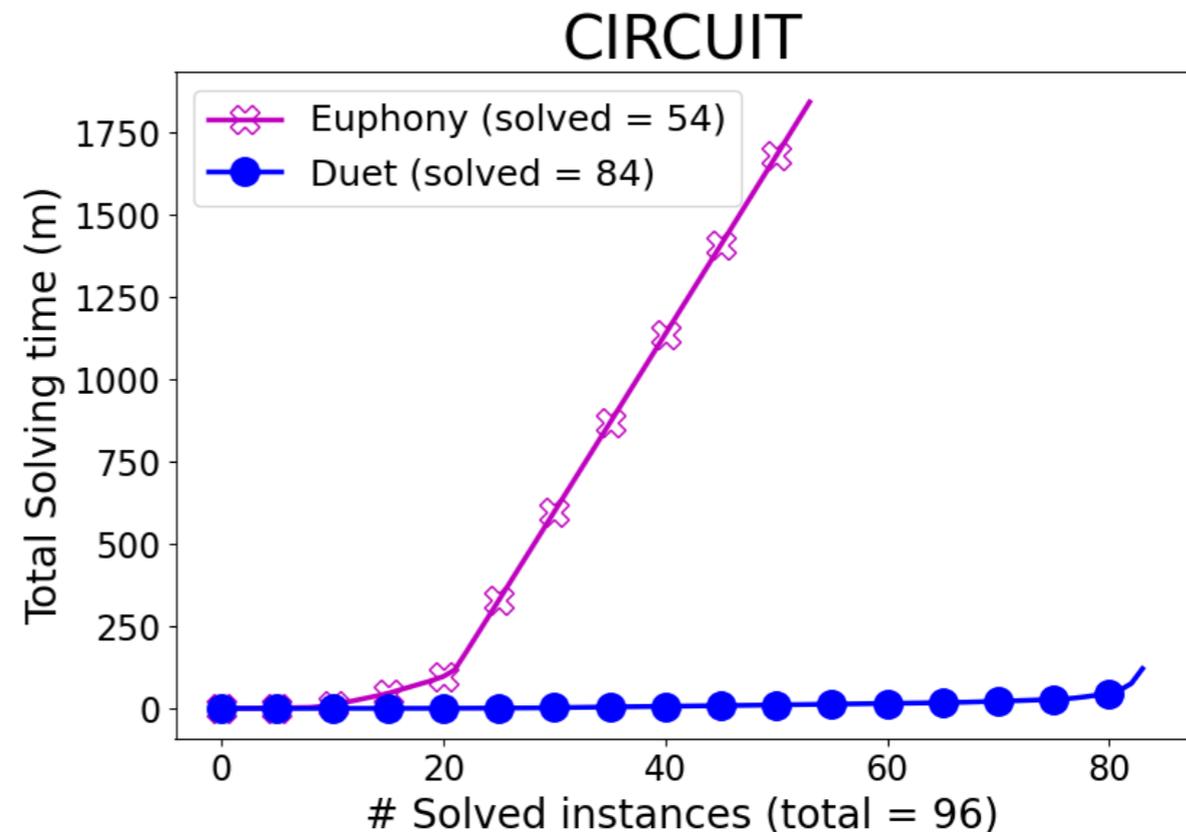
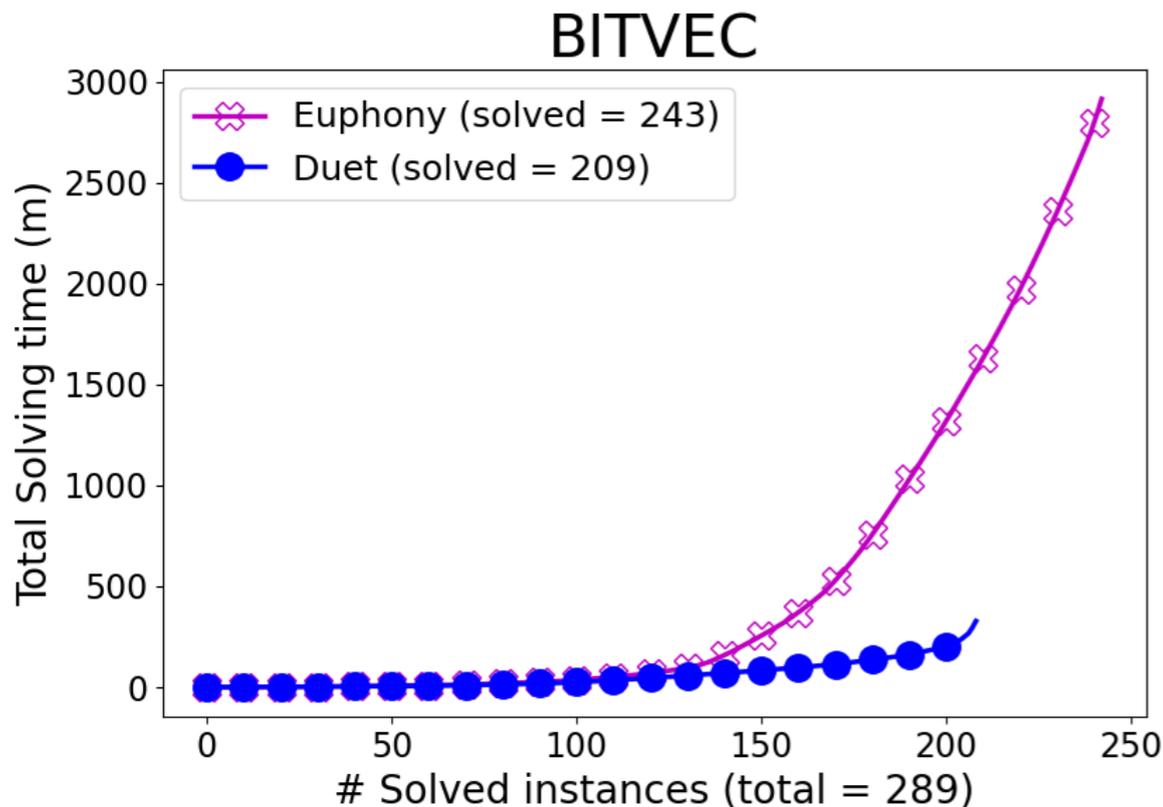
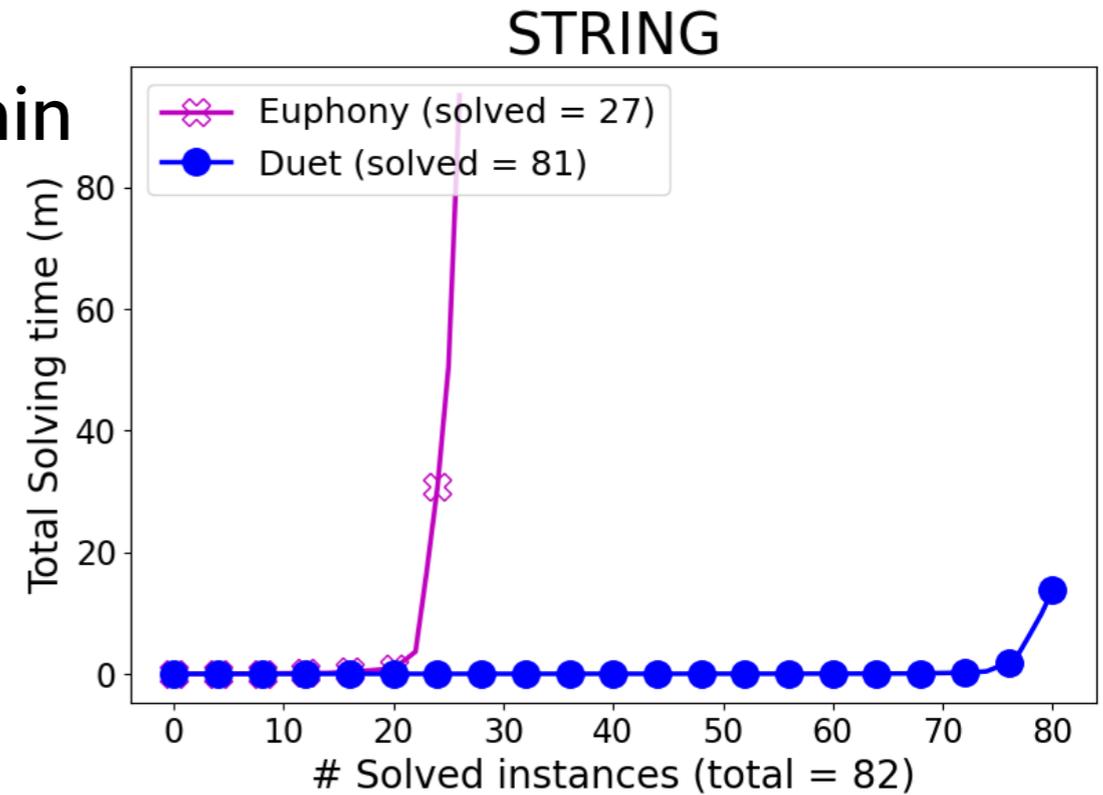
40% within 1 sec

77% within 10 secs

90% within 1 min.

Comparison to Euphony

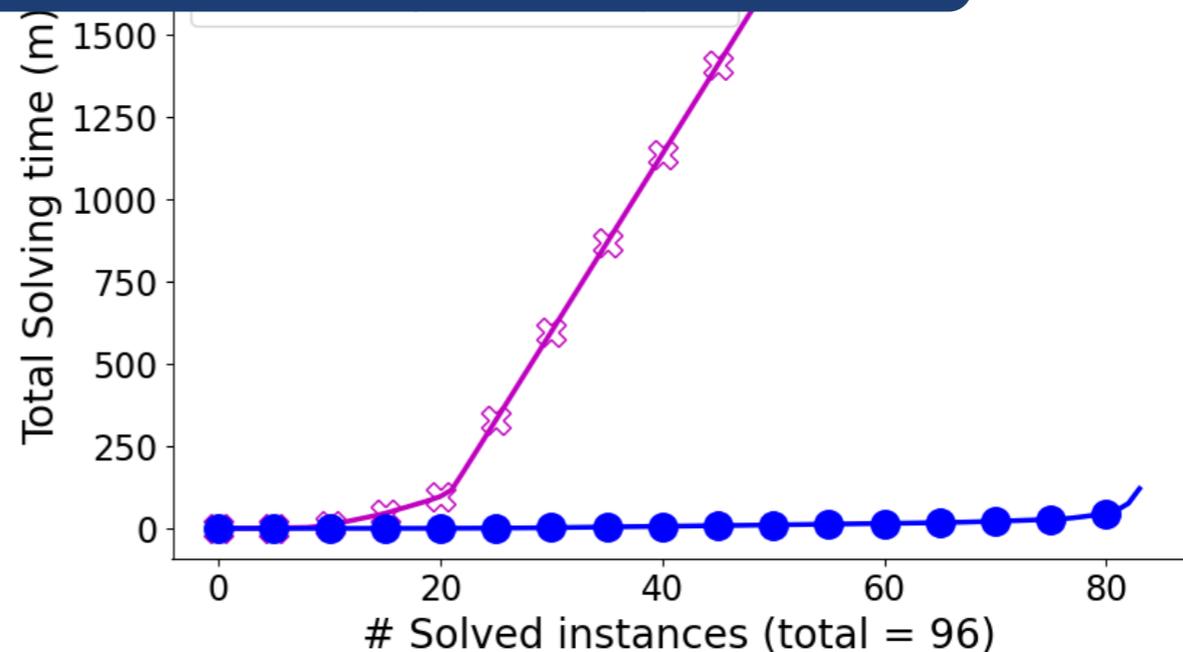
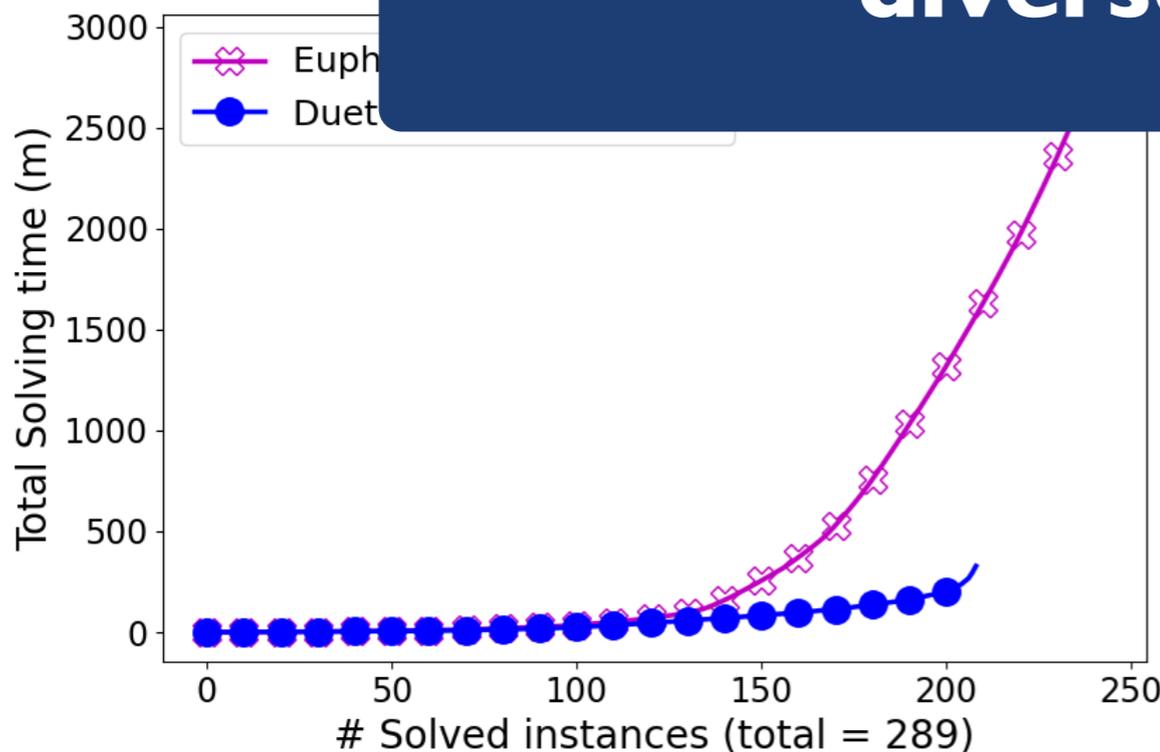
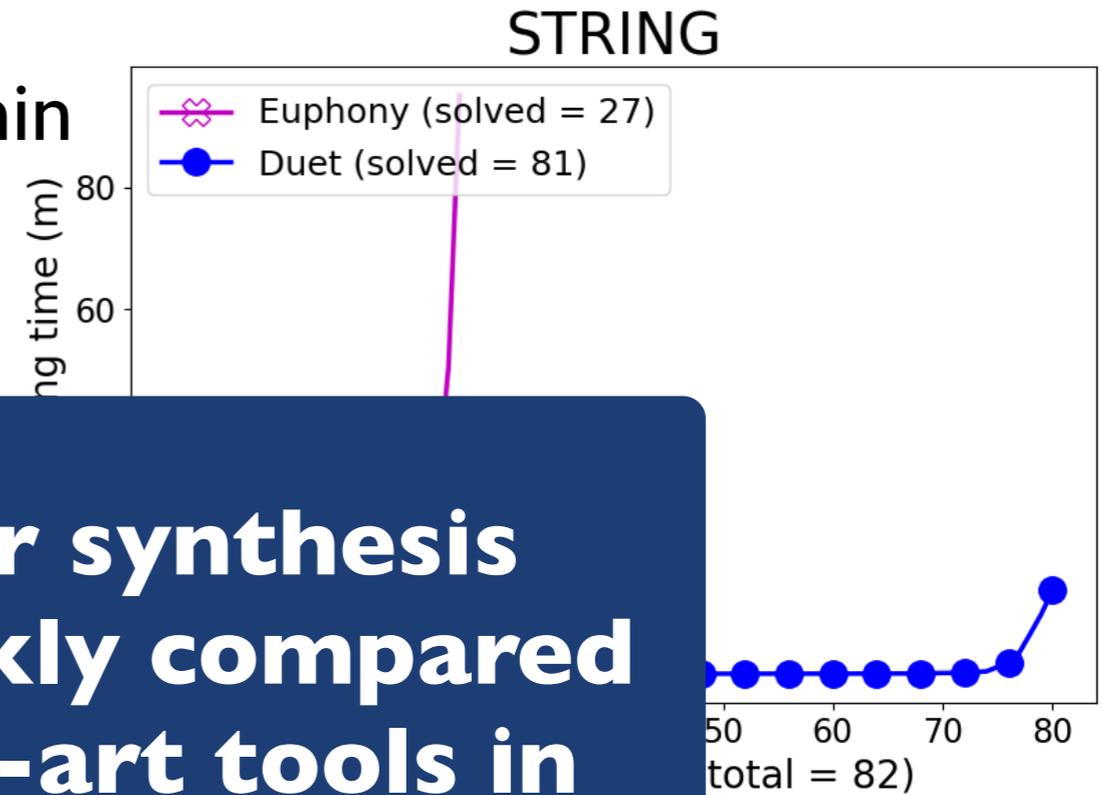
- Training: 1,069 solved EUSolver in 10 min
- Testing: 467
- # solved: Duet: 374, Euphony 324



Comparison to Euphony

- Training: 1,069 solved EUSolver in 10 min
- Testing: 467
- # solved

Duet solves harder synthesis problems more quickly compared to the state-of-the-art tools in diverse domains.



Comparison to Euphony

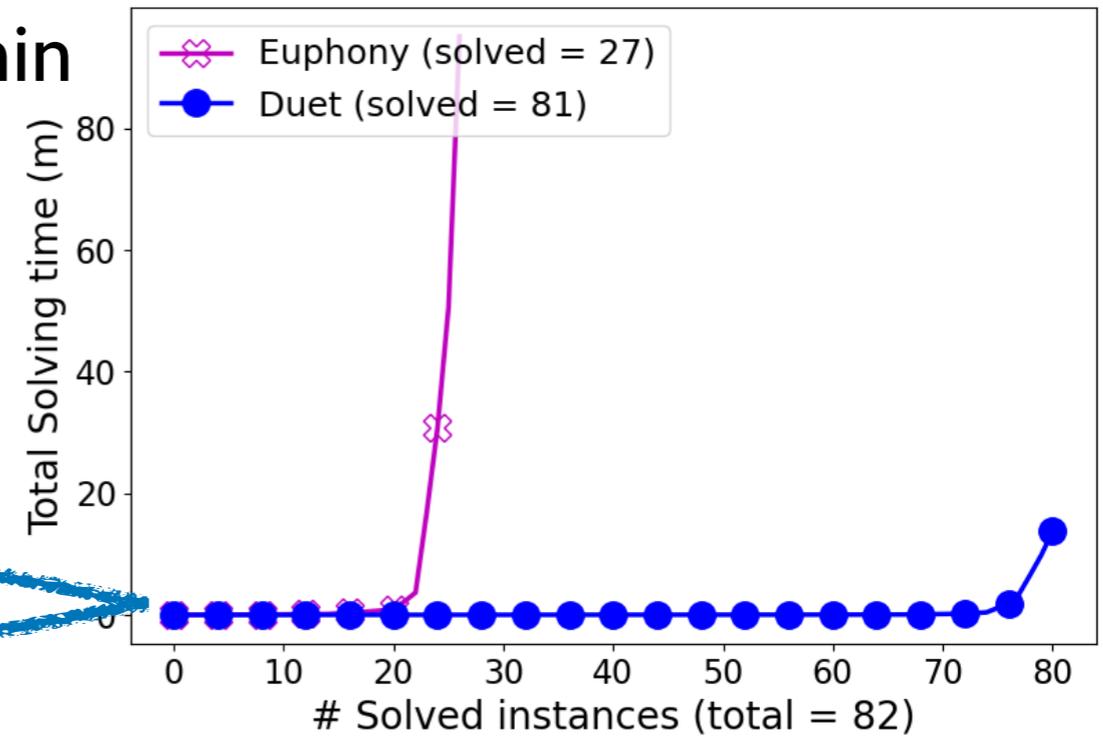
- Training: 1,069 solved EUSolver in 10 min

- Testing: 4

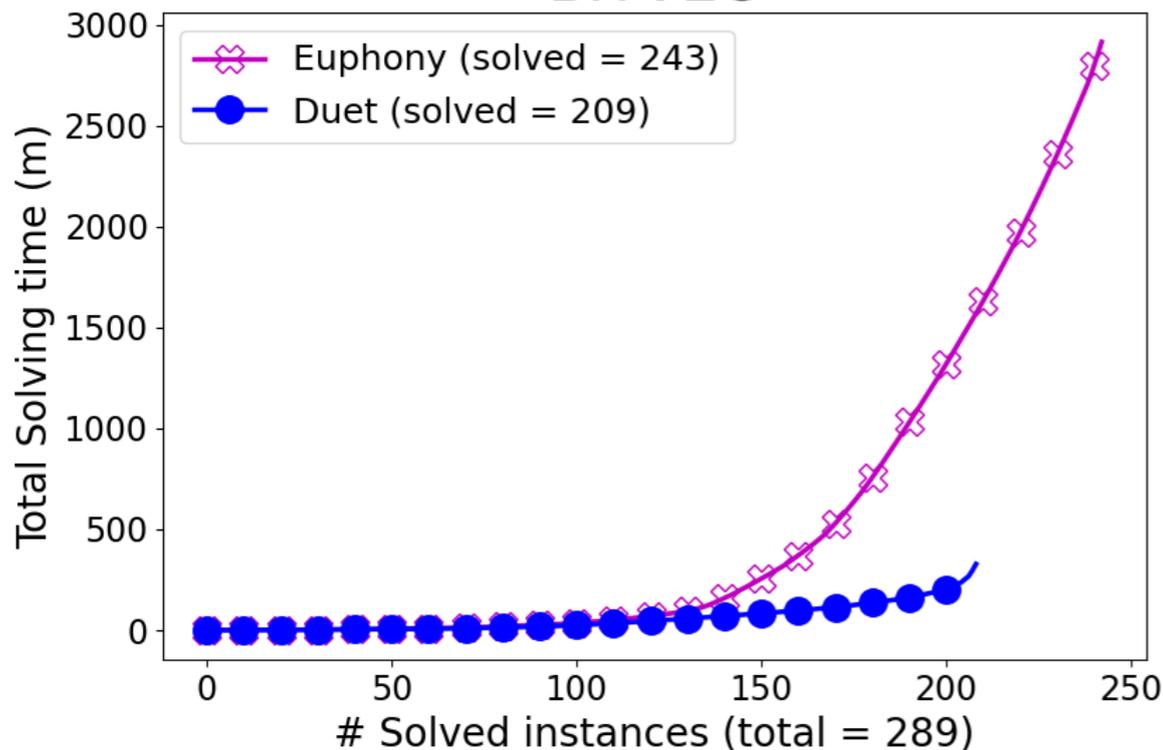
- # solved

Euphony is outperformed by Duet despite the use of statistical models.

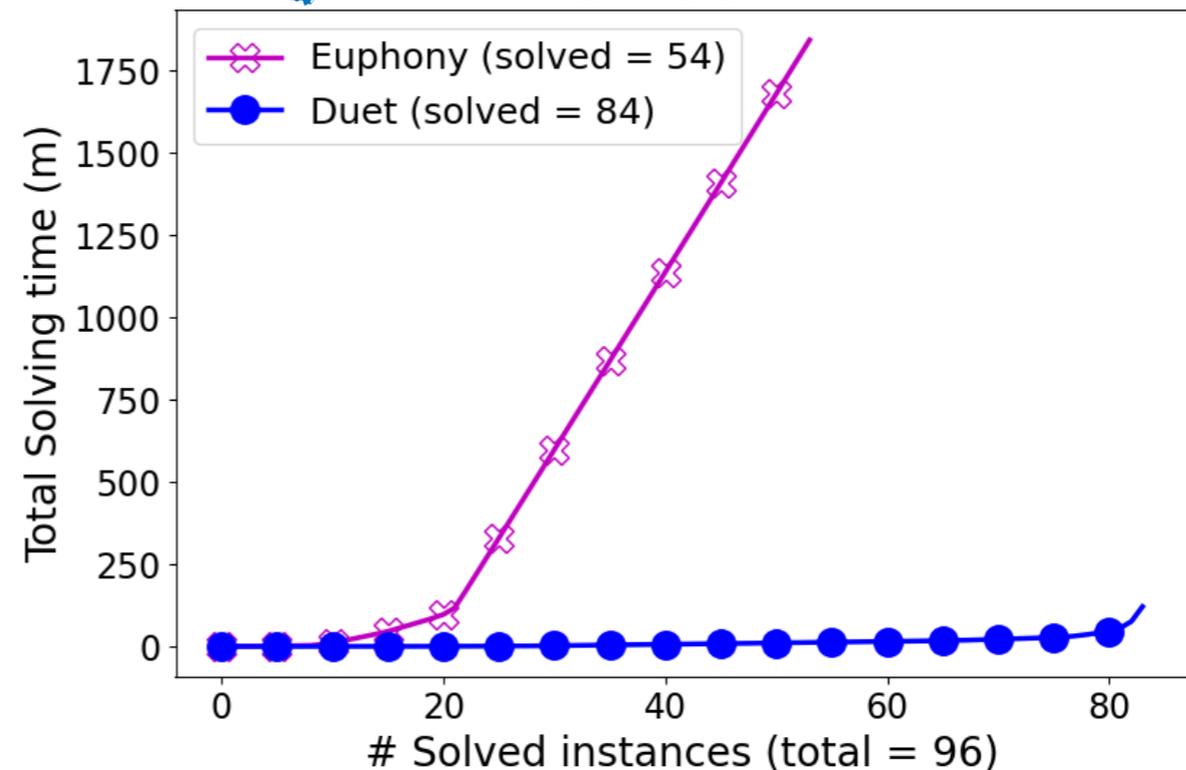
STRING



BITVEC



CIRCUIT

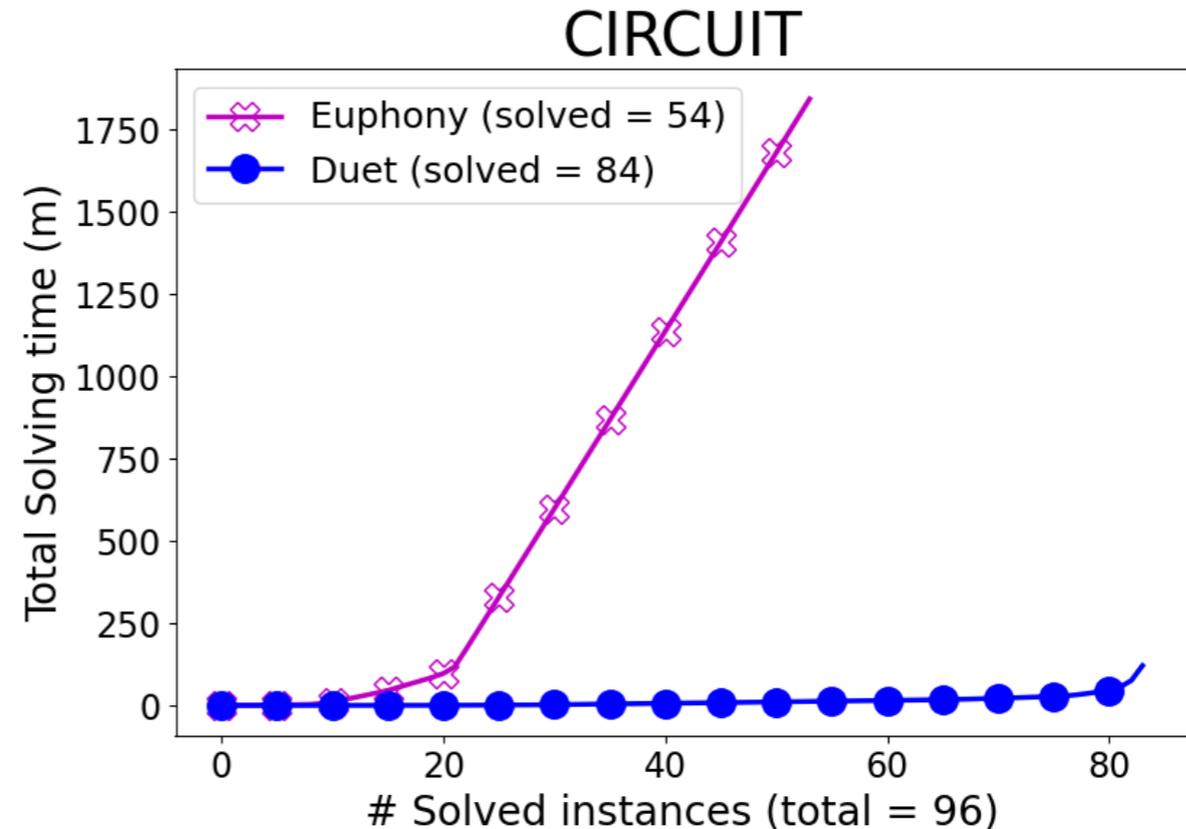
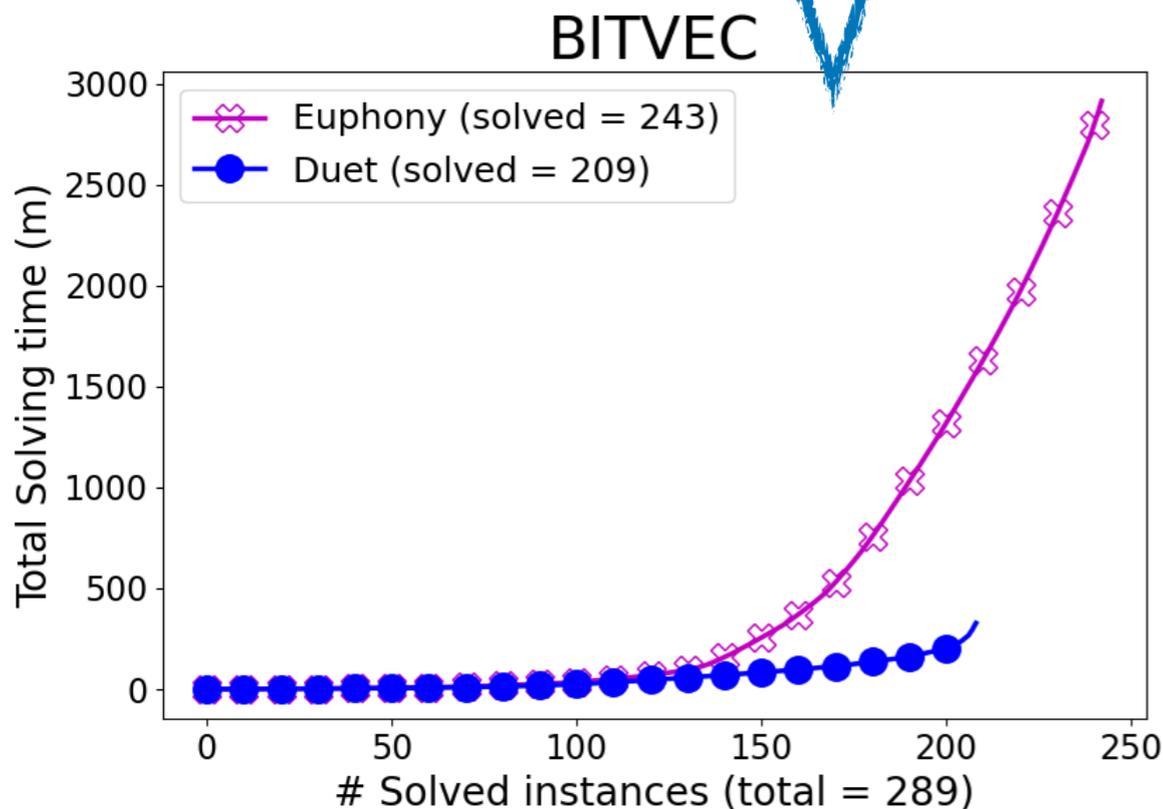
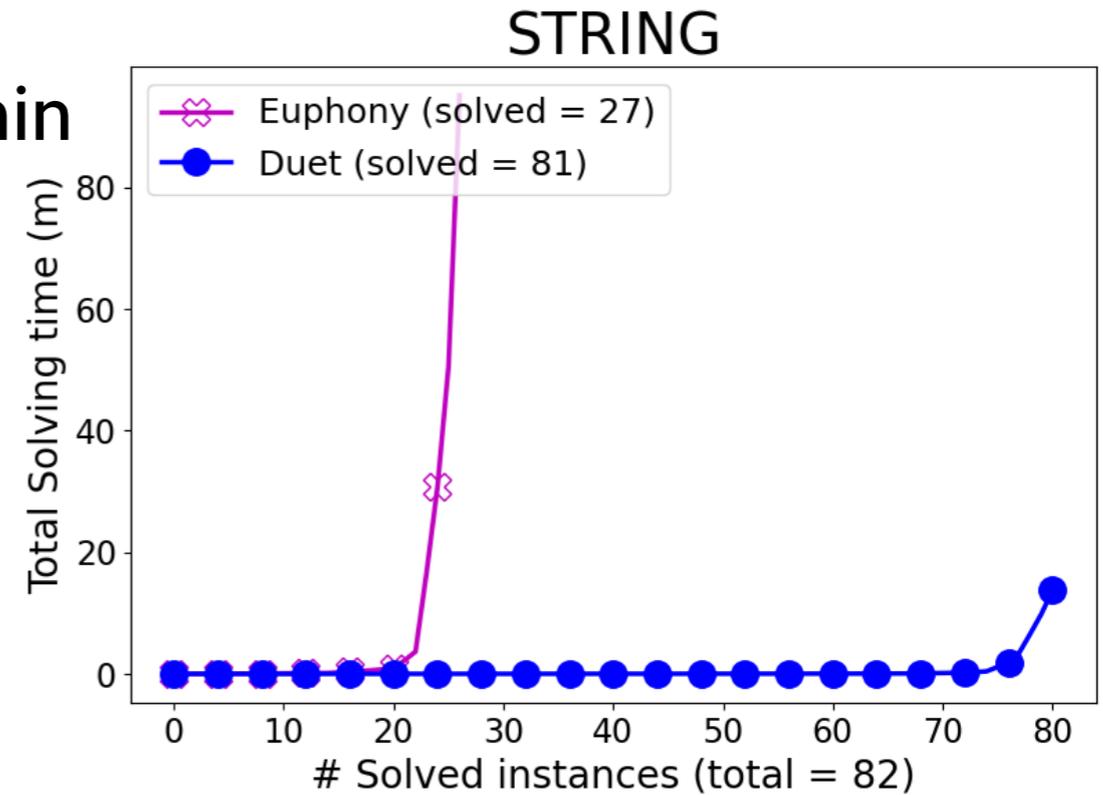


Comparison to Euphony

- Training: 1,069 solved EUSolver in 10 min

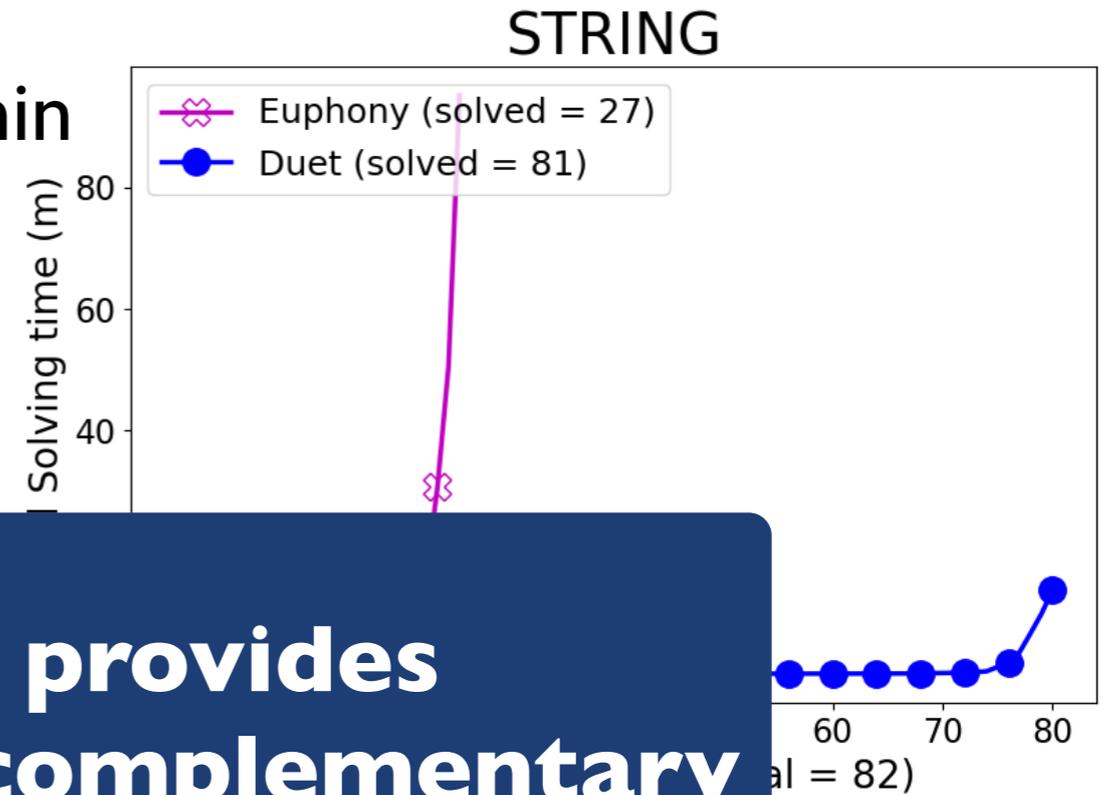
- Testing: 467

- # solved Euphony outperforms Duet by exploiting statistical regularity.

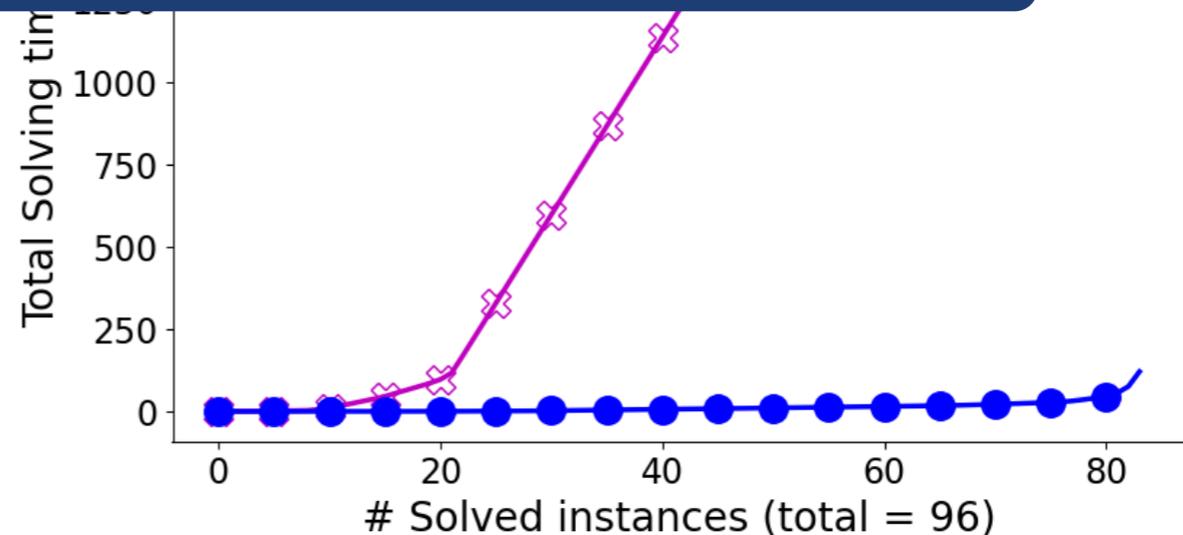
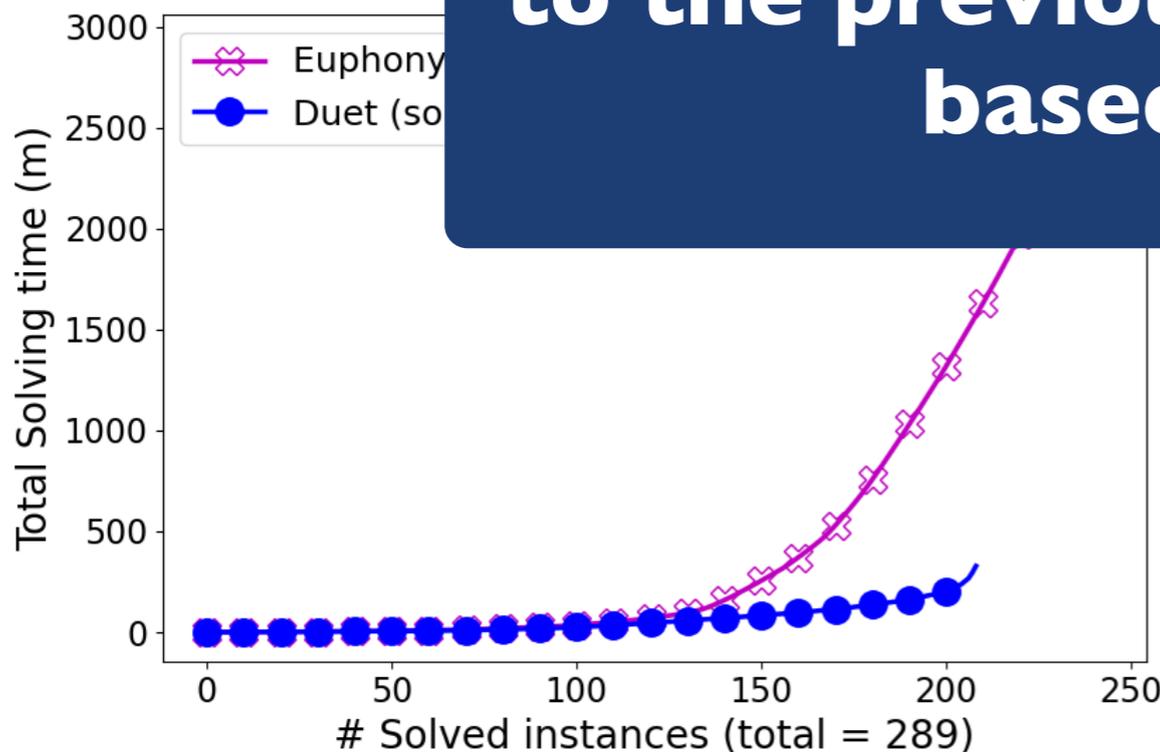


Comparison to Euphony

- Training: 1,069 solved EUSolver in 10 min
- Testing: 467
- # solved: D



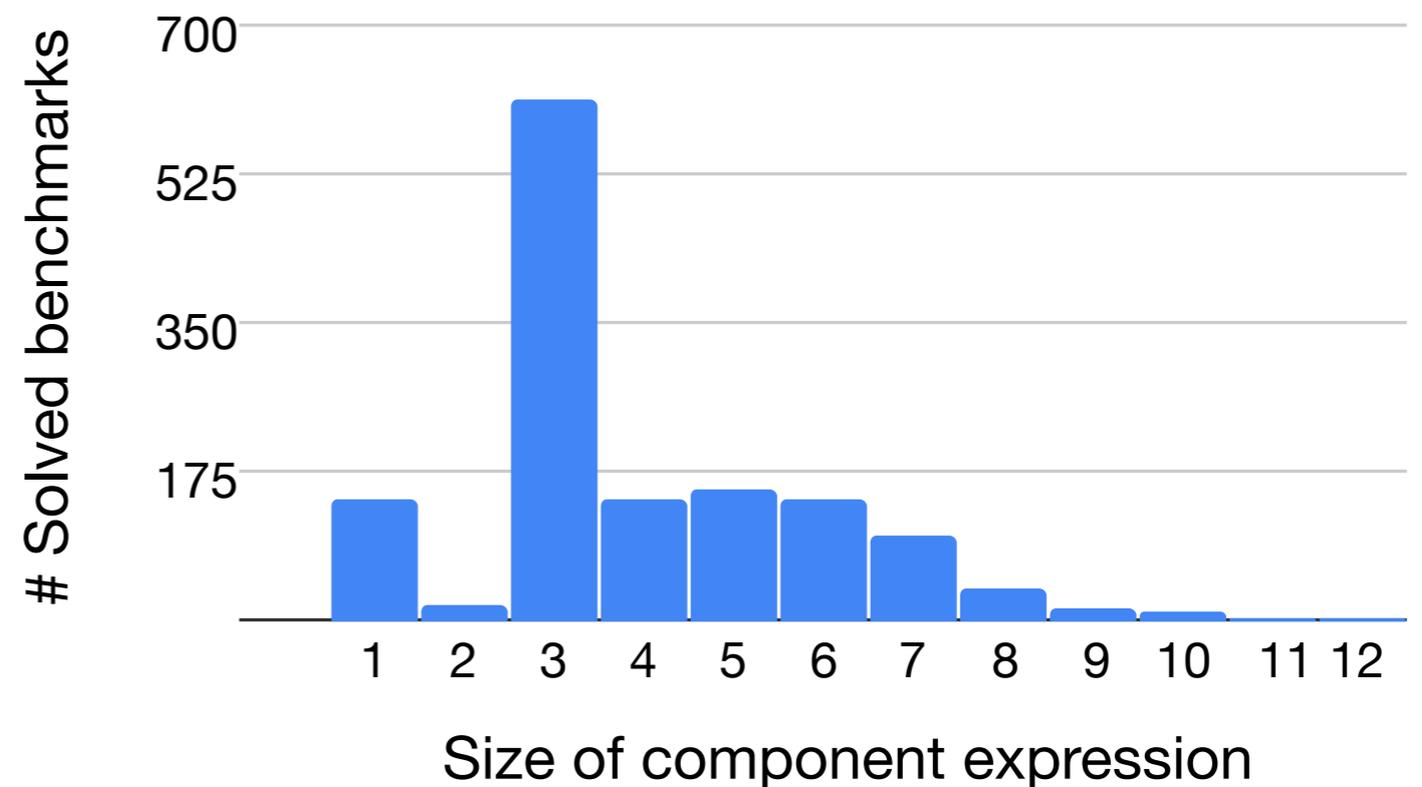
Our approach provides performance gains complementary to the previous machine-learning-based technique.



Analysis of Component Sizes

- Sizes of components needed to construct a solution

- STRING: 1 ~ 6 (avg: 1.8)
- BITVEC: 3 ~ 5 (avg: 3.5)
- CIRCUIT: 3 ~ 12 (avg: 5.9)



- 92% of problems could be solved with components of size ≤ 7 .

Analysis of Component Sizes

- Sizes of components needed to construct a solution

- STRING: 1 ~ 6 (avg: 1.8)

- BITVECTOR (1 ~ 25)

- CIRCULAR



- 92% of problems could be solved with components of size ≤ 7 .