Sound Non-statistical Clustering of Static Analysis Alarms

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Contents

- Problem & Our approach
- Overall result
- Clusterings
- Framework
- Conclusion

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Motivation

- Manual alarm investigation is painful!
- Our commercial tool Sparrow
 - For 3.5 MLOC program
 - Over 1060 alarms are reported.

Our approach

- Cluster similar alarms of the same origin
- Clusters have its own representatives (= dominant alarms).
- Users may inspect only dominant alarms.

gzip-1.2.4

```
void pqdownheap(int k)
{
    int j = 2 * k;
    while(j <= heap_len)
    {
        heap[k] = heap[j];
        k = j;
        j = 2 * j;
    }
    heap[k] = ...;
}</pre>
```

3 buffer-overflow alarms

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gzip-1.2.4

```
void pqdownheap(int k)
{
    int j = 2 * k;
    while(j <= heap_len)
    {
        heap[k] = heap[j];
        k = j;
        j = 2 * j;
    }
    heap[k] = ...;
}</pre>
```

A user identifies heap [j] to be false

gzip-1.2.4

```
void pqdownheap(int k)
{
    int j = 2 * k;
    while(j <= heap_len)
    {
        heap[k] = heap[j];
        k = j;
        j = 2 * j;
    }
    heap[k] = ...;
}</pre>
```

The others are automatically deduced false.

gzip-1.2.4

```
void pqdownheap(int k)
{
    int j = 2 * k;
    while(j <= heap_len)
    {
        heap[k] = heap[j];
        k = j;
        j = 2 * j;
    }
    heap[k] = ...;
}</pre>
```

Users may check only <u>heap[j]</u> instead of all.

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Results: Example (1/2)

```
char invmergerules[8];
char invmergerules_nn[8];
int lookup (char *rule) {
  for (i = 1; invmergerules[i]; i++)
    if (strcasecmp(rule, invmergerules_nn[i] == 0)
      return (i);
}
int rule (struct sketch *s, int rule, int rcount) {
  if (debug)
    printf("%s %d", invmergerules[rule], rcount);
}
int apply (char *rule, struct sketch *sketch) {
  if (code = lookup (rule))
       res = rule (sketch, code, rcount);
```

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Results: Example (2/2)

```
char cboard[64];
char ephash[64];
void MakeMove(int side, int *move) {
  fpiece = cboard[f];
  tpiece = cboard[t];
  if (fpiece == pawn && abs(f-t) == 16) {
    sq = (f + t) / 2;
    HashKey ^= ephash[sq];
  }
}
```

gnuchess-5.05

Results: Overall Effectiveness

Program	LOC	# Alarms	# Alarms after Clustering	% Reduction
nlkain-1.3	831	124	93	25%
polymorph-0.4.0	I,357	25	13	48%
ncompress-4.2.4	2,195	66	30	55%
sbm-0.0.4	2,467	237	125	47%
stripcc-0.2.0	2,555	194	127	35%
barcode-0.96	4,460	435	302	31%
129.compress	5,585	57	29	49%
archimedes-0.7.0	7,569	711	132	81%
man-1.5hl	7,232	276	165	40%
gzip-1.2.4	11,213	385	263	32%
combine-0.3.3	11,472	733	294	60%
gnuchess-5.05	11,629	976	333	66%
bc-1.06	12,830	593	198	67%
coan-4.2.2	22,414	461	291	37%
grep-2.5.1	31,154	115	85	26%
lsh-2.0.4	110,898	616	264	57%
Total	245,861	6,004	2,744	54%

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Clusterings

- On top of the interval-domain-based industrialized commercial tool Sparrow
- Three alarm clustering analyses
 - I. Syntactic clustering
 - 2. Semantic clustering with non-relational analysis
 - 3. Semantic clustering with relational analysis

I. Syntactic Clustering

- while (*optarg && *optarg >= '0' && *optarg <= '9')
 val = *optarg '0';
 optarg++;</pre>
 - Expressions are the same.
 - Variables have the same definition point.

• key idea (alarm dependence)

```
int buffer[10];
buffer[i] = 10; // i = [0, \infty]
j = i / 3; // j = [0, \infty]
foo = buffer[j]; // j = [0, \infty]
```

Two alarms occurred.

• key idea (alarm dependence)

int buffer[10]; buffer[i] = 10; // i = [0, 9] j = i / 3; // j = [0, ∞] foo = buffer[j]; // j = [0, ∞]

assume buffer[i] false

• key idea (alarm dependence)

```
int buffer[10];
buffer[i] = 10; // i = [0, 9]
j = i / 3; // j = [0, 3]
foo = buffer[j]; // j = [0, 3]
```

propagate the refinement

• key idea (alarm dependence)

```
int buffer[10];
buffer[i] = 10; // i = [0, 9]
j = i / 3; // j = [0, 3]
foo = buffer[j]; // j = [0, 3]
```

It kills the other.

• key idea (alarm dependence)

int buffer[10]; buffer[i] = 10; // i = [0, 9] j = i / 3; // j = [0, 3] foo = buffer[j]; // j = [0, 3]

If **buffer[i]** is false, so is the other.

We cluster two alarms.

```
char * p, * str;
for (p = str; *p; p++) // 0 ≤ p.offset
 *p = T0L0WER(*p);
```

```
if (*str == '*') ... // 0 \le str.offset
Two alarms occurred.
```

```
char * p, * str;
for (p = str; *p; p++) // 0 ≤ p.offset < p.size
 *p = T0L0WER(*p);
```

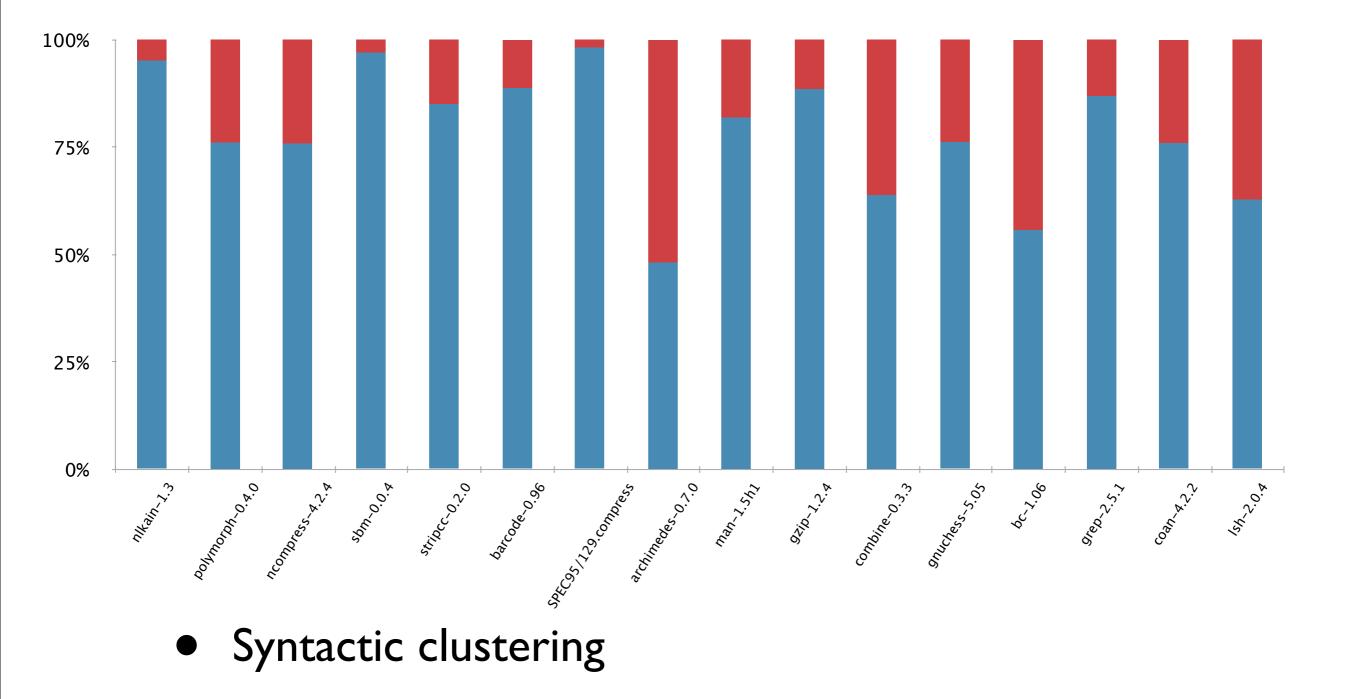
```
if (*str == '*') ... // 0 ≤ str.offset
    assume *p false
```

```
char * p, * str;
for (p = str; *p; p++) // 0 ≤ p.offset < p.size
*p = TOLOWER(*p);
// Loop inv :
// 0 ≤ str.offset ≤ p.offset < p.size = str.size
if (*str == '*') ... // 0 ≤ str.offset < str.size
propagate the refinement
```

```
char * p, * str;
for (p = str; *p; p++) // 0 ≤ p.offset < p.size
*p = TOLOWER(*p);
// Loop inv :
// 0 ≤ str.offset ≤ p.offset < p.size = str.size
if (*str == '*') ... // 0 ≤ str.offset < str.size
It kills the other.
```

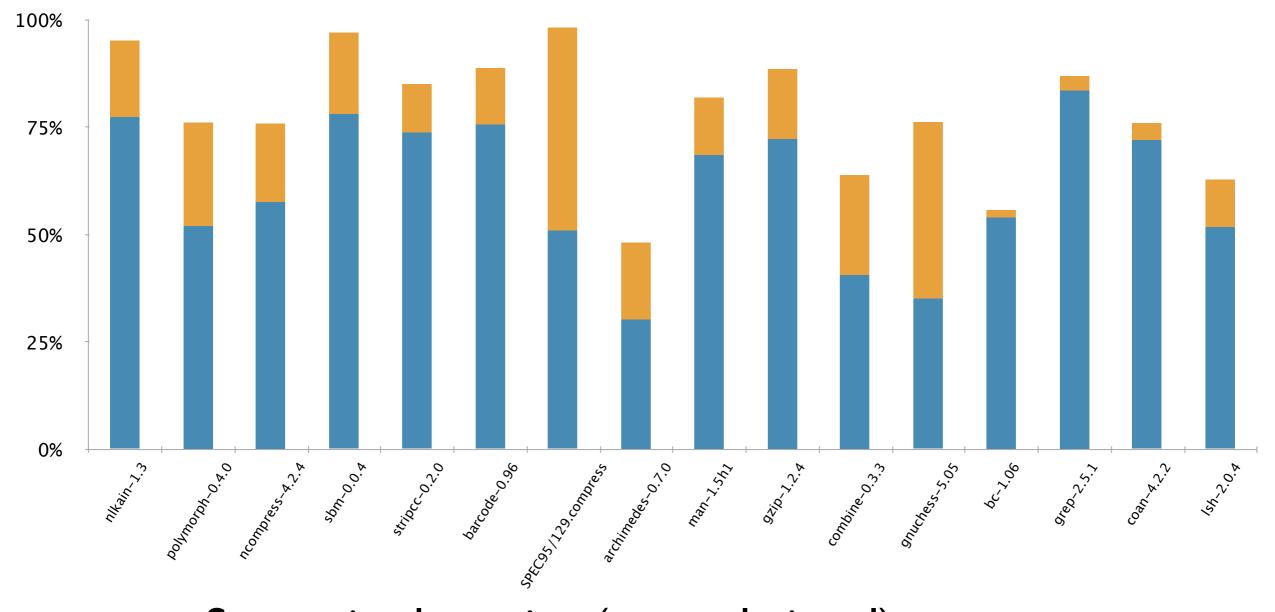
```
char * p, * str;
for (p = str; *p; p++) // 0 ≤ p.offset < p.size
*p = TOLOWER(*p);
// Loop inv :
// 0 ≤ str.offset ≤ p.offset < p.size = str.size
if (*str == '*') ... // 0 ≤ str.offset < str.size
If *p is false, so is the other.
We cluster two alarms.
```

Result



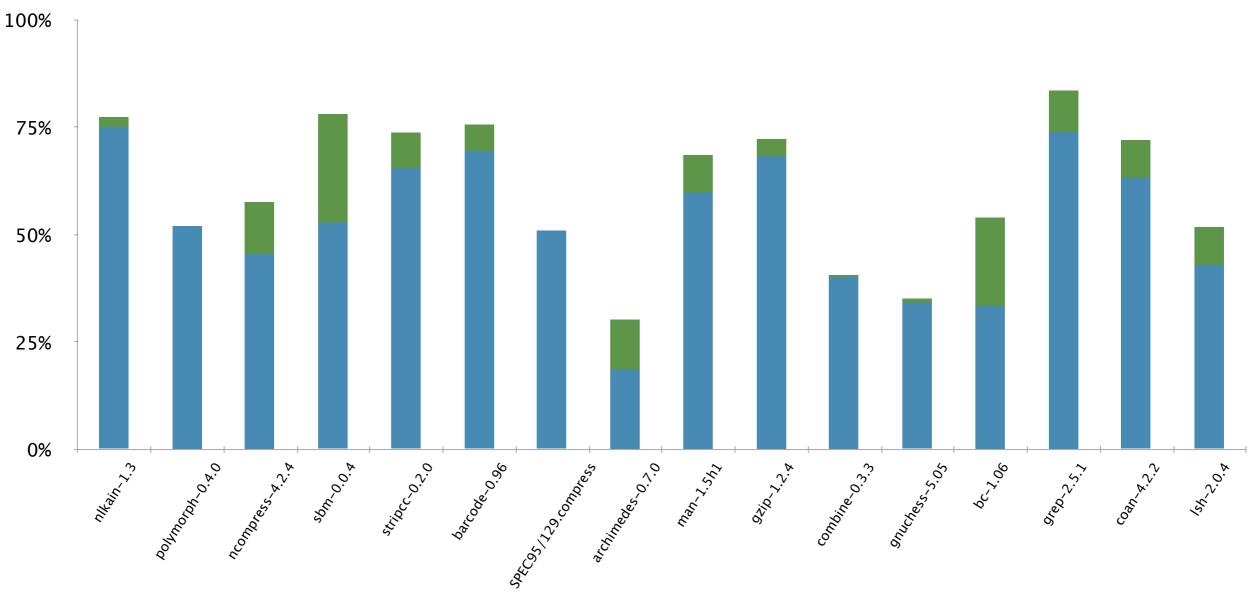
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Result

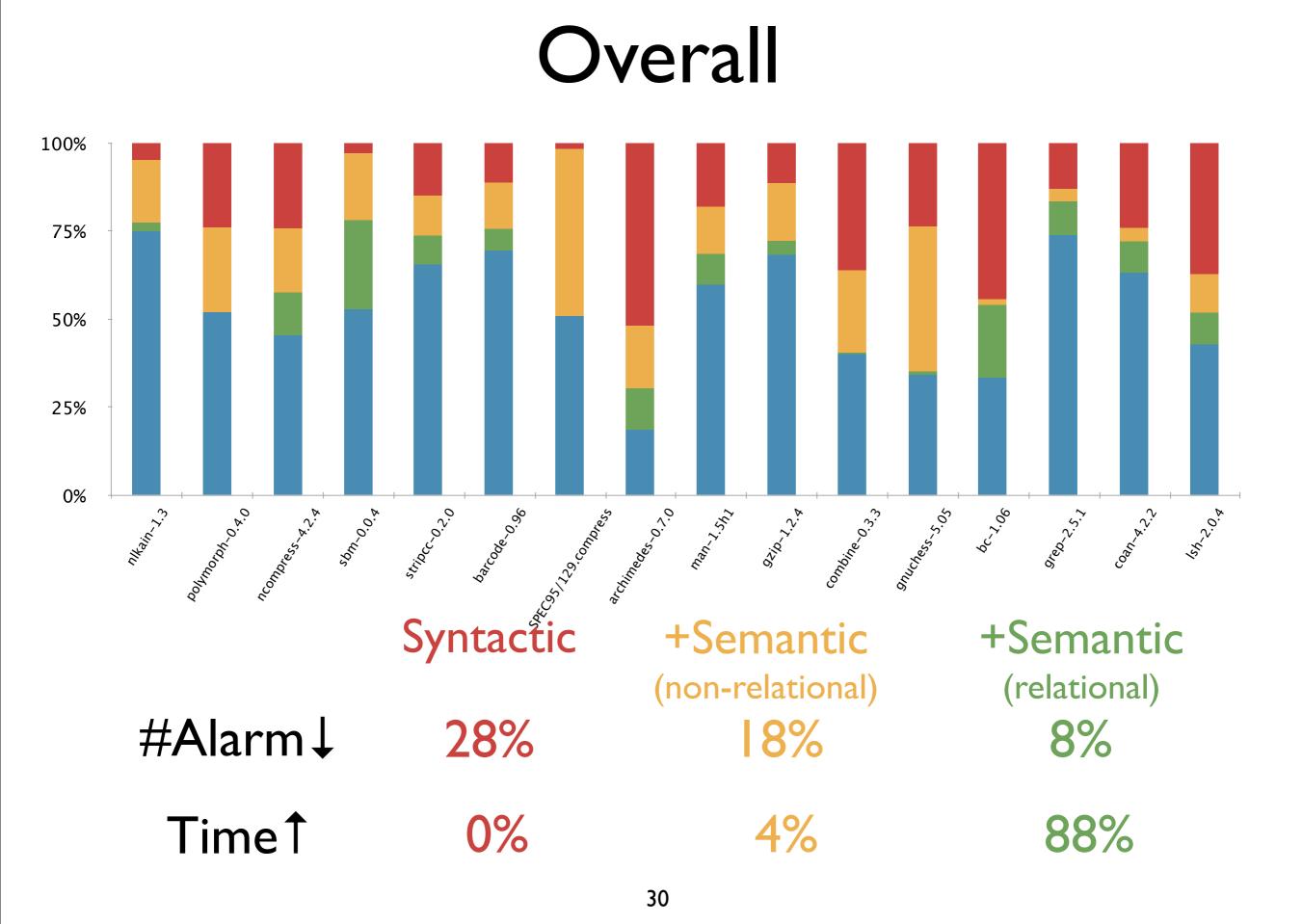


• Semantic clustering (non-relational)

Result



• Semantic clustering (relational)



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Formalization & Soundness

- Three methods have the same strategy.
 - (1) Assume some alarms are false
 (2) propagate the refinement
 (3) get alarm dependences
- We formalize a general alarm clustering method, and prove the correctness.

Alarm Clustering Framework

- Three clusterings are instances of the framework.
- Applicable to any semantics-based static analysis
- Guarantees the soundness of alarm clustering

Notations

- Set of program points $\,\Phi$, and set of the states S
- Concrete semantics $\llbracket P \rrbracket : \Phi \to 2^S$
- Galois connection

$$2^S \xleftarrow{\gamma}{\alpha} \hat{S}$$

• Abstract semantics $\hat{T}: \Phi \to \hat{S}$

$$\forall \varphi \in \Phi.\alpha(\llbracket P \rrbracket(\varphi)) \sqsubseteq \hat{T}(\varphi)$$
$$\hat{T} = \operatorname{fix} \hat{F}$$

• Erroneous states $\Omega: \Phi \to 2^S$

Goal

For any two alarms at φ₁, φ₂ ∈ Φ,
 to find concrete dependence
 [P](φ₁) ∩ Ω(φ₁) = Ø ⇒ [P](φ₂) ∩ Ω(φ₂) = Ø

• Using abstract dependence

Abstract Alarm Dependence $\varphi_1 \rightsquigarrow \varphi_2$

Definition 1 ($\varphi_1 \rightsquigarrow \varphi_2$) Given two alarms φ_1 and φ_2 , alarm φ_2 has abstract dependence on alarm φ_1 if and only if,

 $\gamma(\tilde{T}_{\varphi_1}(\varphi_2)) \cap \Omega(\varphi_2) = \emptyset$

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$$\gamma(\tilde{T}_{\varphi_1}(\varphi_2)) \cap \Omega(\varphi_2) = \emptyset$$

where

$$\tilde{T}_{\varphi_1} = \operatorname{fix} \lambda X. \qquad \hat{F}(X) \qquad (\hat{T} = \operatorname{fix} \hat{F})$$

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$$\hat{T}_{\neg\varphi_1} = \hat{T}\{\varphi_1 \mapsto \hat{T}(\varphi_1) \,\hat{\ominus} \,\alpha(\Omega(\varphi_1))\}$$

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$$\hat{T}_{\neg\varphi_1} = \hat{T}\{\varphi_1 \mapsto \underline{\hat{T}(\varphi_1) \ominus \alpha(\Omega(\varphi_1))}\}$$

slice out error states at φ_1 in a way that it approximates $\llbracket P \rrbracket(\varphi_1) \ominus \Omega(\varphi_1)$

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$$\gamma(\tilde{T}_{\varphi_1}(\varphi_2)) \cap \Omega(\varphi_2) = \emptyset$$

where

$$\tilde{T}_{\varphi_1} = \underline{\operatorname{fix} \lambda X} \hat{T}_{\neg \varphi_1} \sqcap \hat{F}(X)$$

propagate the refinement until fixpoint

$$\hat{T}_{\neg\varphi_1} = \hat{T}\{\varphi_1 \mapsto \underline{\hat{T}(\varphi_1) \ominus \alpha(\Omega(\varphi_1))}\}$$

slice out error states at $\,arphi_1$ in a way that

it approximates

 $\llbracket P \rrbracket(\varphi_1) \ominus \Omega(\varphi_1)$

Soundness of $\varphi_1 \rightsquigarrow \varphi_2$

Lemma 1 $\varphi_1 \rightsquigarrow \varphi_2 \implies (\text{alarm } \varphi_1 \text{ false} \implies \text{alarm } \varphi_2 \text{ false})$

φ_1 can be lifted to a set of alarms

Lemma 2 $\overrightarrow{\varphi} \rightsquigarrow \varphi \Longrightarrow ((\forall \varphi_i \in \overrightarrow{\varphi}.alarm \ \varphi_i \ false) \Longrightarrow alarm \ \varphi \ false)$

Alarm Cluster $C_{\vec{\varphi}}$

Definition 3 (Alarm Cluster) Given set \mathcal{A} of all alarms and dependence relation \rightsquigarrow , a false alarm cluster $\mathcal{C}_{\overrightarrow{\varphi}}$ is $\{\varphi \in \mathcal{A} \mid \overrightarrow{\varphi} \rightsquigarrow \varphi\}$.

 $\overrightarrow{\varphi} \subseteq \mathcal{A}$: dominant alarms of cluster $\mathcal{C}_{\overrightarrow{\varphi}}$

Clustering Algorithm

- Dependences determine the clustering.
- Brute-force search requires $2^{\text{#Alarms}}$ fixpoint computation.
- Our algorithm requires one fixpoint computation.
 - but misses some dependences.
- <u>The algorithm derives sound dependences.</u>[†]

[†]Not in the paper. Please refer to technical memo : <u>http://ropas.snu.ac.kr/~wslee/vmcai12_techmemo.pdf</u>

Conclusion

A sound, general, and effective way to reduce alarm-investigation efforts

Thank you!

Backup slides

Example

	<pre>int large[7]; int medium[5]; int small[3];</pre>	$ ilde{T}_{\mathcal{A}}$	R
$arphi_1$	large[i] =;	$\dots [0, 6]$	$\{\varphi_1\}$
$arphi_2$	= medium[i]; ····	$\cdots [0,4]$	$\{\varphi_2\}$
$arphi_3$	= large[i];	$\cdots [0, 4]$	$\{\varphi_2\}$
$arphi_4$	= medium[i-1];	$\cdots [1, 4]$	$\{ arphi_2, \ arphi_4 \}$
$arphi_5$	= small[i-1];	[1, 4]	$\{ \varphi_2, \ \varphi_4 \}$

• Clustering result $C_{\varphi_2} = \{\varphi_3\}$

$$C_{\{\varphi_2,\varphi_4\}} = \{\varphi_5\}$$

- Naive algorithm
 - $C_{\varphi_1} = \{\varphi_3\}$ $C_{\varphi_2} = \{\varphi_3\}$

 $C_{\{\varphi_2,\varphi_4\}} = \{\varphi_5\}$

Experimental result

 Table 1. Alarm clustering results.

 ${\bf B}$: baseline analysis, ${\bf S}$: syntactic alarm clustering, ${\bf I}$: semantic alarm clustering with interval domain, ${\bf O}$: semantic clustering with octagon domain.

Program	LOC	# Alarms			% Reduction			Time(s)				
		B	S	S+I	S+I+O	S	+I	+0	S+I+O	B	Ι	0
nlkain-1.3	831	124	118	96	93	5%	18%	2%	25%	0.17	0.03	0.1
polymorph-0.4.0	1,357	25	19	13	13	24%	24%	0%	48%	0.12	0	0.06
ncompress-4.2.4	2,195	66	50	38	30	24%	18%	12%	55%	0.54	0.03	0.69
sbm-0.0.4	2,467	237	230	185	125	3%	19%	25%	47%	2.28	0.3	1.15
stripcc-0.2.0	2,555	194	165	143	127	15%	11%	8%	35%	2.76	0.07	25.44
barcode-0.96	4,460	435	386	329	302	11%	13%	6%	31%	3.23	0.1	2.59
129.compress	5,585	57	56	29	29	2%	47%	0%	49%	2.46	0.02	0.19
archimedes-0.7.0	7,569	711	342	215	132	52%	18%	12%	81%	6.48	0.27	16.11
man-1.5h1	7,232	276	226	189	165	18%	13%	9%	40%	11.65	0.28	1.86
gzip-1.2.4	11,213	385	341	278	263	11%	16%	4%	32%	10.03	0.3	2.92
combine-0.3.3	11,472	733	468	297	294	36%	23%	0%	60%	19.74	0.81	26.93
gnuchess-5.05	11,629	976	744	343	333	24%	41%	1%	66%	42.49	4.78	8.66
bc-1.06	12,830	593	330	320	198	44%	2%	21%	67%	33.75	7.04	27.23
grep-2.5.1	31,154	115	100	96	85	13%	3%	10%	26%	4.19	0.01	11
coan-4.2.2	22,414	461	350	332	291	24%	4%	9%	37%	126.66	1.91	6.14
lsh-2.0.4	110,898	616	387	319	264	37%	11%	9%	57%	115.13	2.12	204.12
TOTAL	$245,\!861$	6,004	4,312	3,222	2,744	28%	18%	8%	54%	381.68	15.94	335.19